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## Cryptographic Techniques Overview

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| <b>1. Name of Cryptographic Technique</b><br>EPOC   |  |
| <b>Categories</b>   | 1.Asymmetric Cryptographic Schemes<br>2.Symmetric Ciphers<br>3.Hush Functions<br>4.Pseudo-random Number Generators |
| <b>Security Functions of Asymmetric Cryptographic Schemes</b><br>1.confidentiality    2. Authentication    3. signature    4. key- sharing  |  |
| <b>Subcategories of Symmetric Ciphers</b><br>1. stream ciphers    2. 64-bits block ciphers    3. 128-bits block ciphers   |  |
| <b>2. Cryptographic Techniques Overview</b>   |  |
| <p><b>2.1 Design policy</b></p> <p>The design target of EPOC is as follows:</p> <ol style="list-style-type: none"> <li>(1) It should be proven to be secure in the strongest sense (i.e., semantically secure against adaptively chosen-ciphertext attacks) under reasonable assumptions (and in the random oracle model).</li> <li>(2) Its performance should be comparable to the RSA and other practical encryption schemes based on the factoring assumption.</li> <li>(3) Its hybrid usage with a symmetric encryption should be also proven to be secure in the strongest sense (i.e., semantically secure against adaptively chosen-ciphertext attacks) under reasonable assumptions (and in the random oracle model).</li> </ol> <p>Our approach to construct EPOC is based on the random oracle model, in which a primitive public-key encryption function is converted to an encryption scheme provably secure in the strongest sense if the underlying hash functions are assumed truly random functions.</p> <p>Our primitive encryption function is the OU (Okamoto-Uchiyama) function [2], in which to invert the OU function is proven to be as hard as factoring a composite integer public-key). There are three conversions based on the random oracle model [3,4,5], therefore we have three versions of EPOC: EPOC-1, EPOC-2 and EPOC-3.</p> <p>These schemes satisfy the above-mentioned target (security, performance and hybrid security). (EPOC-2 and EPOC-3 for the hybrid security)</p> |  |
| <p><b>2.2 Intended applications</b></p> <ol style="list-style-type: none"> <li>(1) EPOC-1:             <ul style="list-style-type: none"> <li>- Key distribution for a symmetric encryption (at most 256 bit key size)</li> <li>- Encrypted communication for small size data (at most 256 bit data size)</li> </ul> </li> <li>(2) EPOC-2:             <ul style="list-style-type: none"> <li>- Key distribution for a symmetric encryption (no restriction on the size)</li> <li>- Encrypted communication in a hybrid usage with symmetric encryption, especially envelope type (key distribution and data transmission are synchronized)</li> </ul> </li> <li>(3) EPOC-3:             <ul style="list-style-type: none"> <li>- Key distribution for a symmetric encryption (no restriction on the size)</li> <li>- Encrypted communication in a hybrid usage with symmetric encryption, especially "envelope type" (key distribution and data transmission are synchronized)</li> <li>- Encrypted communication in a hybrid usage with symmetric encryption, especially "session type" (only once key distribution in the opening phase of a session, and many times data transmissions during the session)</li> </ul> </li> </ol>   |  |



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### 2.3 Basic theory and techniques

- (1) The elliptic curve ElGamal encryption function as a primitive encryption function.
  - (2) Our novel three conversion methods [2,3,4], by which we have three versions: PSEC-1, PSEC-2, PSEC-3. Especially PSEC-2 and PSEC-3 are the first public-key encryption schemes whose hybrid usages with symmetric encryption are proven to be secure in the strongest sense under reasonable assumptions and random oracle model.
  - (3) In the conversion of PSEC-3 [4], "session type" (only once key distribution in the opening phase of a session, and many times data transmissions during the session) of a hybrid usage with symmetric encryption is available.
- In addition, the overhead of the conversion is almost nothing if practical hash functions such as SHA-1 are employed, namely the conversion is optimal in the performance.

#### References:

- [1] Bellare, M. and Rogaway, P. : Optimal Asymmetric Encryption, Proc. of Eurocrypt'94, LNCS 950, Springer-Verlag, pp.92-111 (1995).
- [2] Fujisaki, E. and Okamoto, T.: How to Enhance the Security of Public-Key Encryption at Minimum Cost, Proc. of PKC'99, Springer-Verlag, LNCS 1560, pp. 53--68 (1999).
- [3] Fujisaki, E. and Okamoto, T.: Secure Integration of Asymmetric and Symmetric Encryption Schemes, Proc. of Crypto'99, Springer-Verlag, LNCS 1666, pp. 535--554 (1999).
- [4] Okamoto, T. and Pointcheval, D.: OCAC: an Optimal Conversion for Asymmetric Cryptosystems, manuscript (2000).