

PDP-1 COMPUTER
ELECTRICAL ENGINEERING DEPARTMENT
M.I.T.
CAMBRIDGE, MASSACHUSETTS 02139

PDP-23-1

ID

May 20, 1966

ID - Invisible Debugger

Invisible Debugger, commonly referred to as ID, is a utility program in the PDP-1 time sharing system written to aid in the debugging of other programs. An advanced ID has been written (April, 1966) to allow all operations to be carried out either directly on drum fields or on running cores. It uses the drum to allow the user full use of core(s) and drum field(s) for his program and to provide extra features. ID and the program being debugged each have a drum field to themselves.

For clarity when typing examples are given herein, the typing done by the user of ID is underlined. Also, when needed the following symbols are assigned to the invisible flexo characters:

carriage return	↵
tabulation	→
space	⊠
backspace	←
upper case	↑
lower case	↓

A. General Essential Preparation

1. When a time-sharing user requests ID he is automatically assigned one drum field to be used for his ID program. The user's running field, which was assigned when the console was turned on, and the console's pseudo drum fields will be used as the console's drum and core fields whose contents may be examined and modified by the use of ID. The drum field assigned upon requesting ID is the console's ID field and may not be examined or modified by instructions to ID or by execution of a program.

2. When entering ID, the user has either a binary tape containing the program and its symbols or a binary version of his program existing on his pseudo field 1 with its symbols still in the POSSIBLE SYMBOL TABLE located in the console's running field (i.e., in core 0).
 - a. Program on Field 1 and Symbols in POSSIBLE SYMBOL TABLE:
To inform ID of the meaning of the symbols used in the program type:

2↑ T ↓ ↷

ID will then take a copy of POSSIBLE SYMBOL TABLE and put it into its own ID SYMBOL TABLE. To get a copy of the binary program from pseudo field 1 and place it into the console's current field so that it can be executed, type:

1↑ U ↓ ↷

(NOTE: If changed, the limits M+1 and M+2 should be initialized before the above command by typing "M".)

- b. Program and Symbols on Tape: To clear all available registers of the current field memory, type:

↑Z↓ ↷

This command will zero all the registers of the current field. Then to kill the previous symbol table, leaving only the initial PDP-1 instruction mnemonics, type:

↑K↓ ↷

To read in the binary tape containing the program or data, place the tape in the reader and type:

↑Y↓ ↷

This causes ID to yank a standard binary block format tape into the current field memory. To inform ID of the meaning of the symbols used in the program, place the symbol tape, which was prepared by POSSIBLE or MIDAS SYMBOL PUNCH, in the reader. If the tape is a binary tape from MIDAS type:

1↑T↓ ↷

If it is an alphanumeric tape from POSSIBLE type:

↑T↓ ↷

ID will then read in the symbol tape and will merge the contents of this tape with ID's own symbol table. After this, ID is ready for use and will be able to interpret constants and instructions typed either symbolically or numerically or both.

3. Typing

↑G↓ ↷

preceded by the address where the user wants his program to begin, will cause the program to start running.

For example, typing:

100↑G↓

or

a↑G↓ ↷

will cause control to be transferred to absolute location 100 or symbolic location "a" respectively.

4. To return control to ID after using a "G" command, press the console's CALL BUTTON.

B. The Current Field

This new version of ID (April, 1966) allows operations to be carried out directly on drum fields or on the user's running cores by making the field involved the "current field". Initially ID is set up so that the user's running core 0 is the "current field". The "current field" is normally specified by the underbar command. Typing

x_

causes field x to become the "current field". If $x \leq 177$, x itself is used; otherwise, bits 2-5 of x are used. Fields 0 to 7 refer to the user's normal running core. For time-sharing users, only 0 and 1 are legal, and 1 is legal only if core 1 is assigned to the user. Fields 10-100 are illegal. Fields greater than (>) 100 refer to drum fields. For example, typing

103_

makes the user's pseudo field 3 the current field. Absolute fields are indicated by bit 12. For example,

161_

causes absolute field 21 to become "current". Typing underbar (_) alone will cause the current field to be printed out.

C. Examination and Modification of Stored Information

1. Opening a register in the current field - In using ID, a fundamental idea is that of opening a register so that its contents may be examined and/or changed. This may be accomplished by typing the twelve-bit address of the register in the current field to be opened, either symbolically or as an absolute constant, followed by a slash. For example:

reg+2/

or

2467/

When the above is typed ID will immediately print a tabulation, then the contents of that register in the current field, followed by another tabulation. Continuing the example above:

reg+2/ →| add loc+3 →|

(NOTE: Current drum fields not assigned to the user cannot be examined.)

2. Examination of a register not in the current field - It is frequently desirable to open a register not in the current field so that its contents may be examined and/or changed. This is accomplished by typing a 16 bit extended address of the register to be opened, either symbolically or as an absolute constant, followed by a vertical bar. The corresponding core module will become the "current field". For example:

12345|

will cause core 1 to become current, and open register 2345. Typing

1 reg|

will open register reg in core 1 and make core 1 current.

(NOTE: 1=10000 in POSSIBLE and ID symbol tables.)

Like slash, when the above is typed, ID will immediately print a lower case, tabulation, then the contents of the register, followed by another tabulation.

(NOTE: For time-sharing users, reference to core 1 is legal only if core 1 is assigned to that user.)

3. Modifying and closing a register - Once a register has been opened in either of the above manners its contents may be modified, if desired, by typing the change either symbolically or as a constant. For example:

reg+2/ →| add loc+3 →| add loc+5

(NOTE: System fields or drum fields not assigned to the user cannot be modified.)

A command character which may be helpful in modifying registers is Q. This has the value of the last quantity typed by ID or you. For example, to change the contents of register 50 from 155 to 157 type:

50/ →| 155 →| Q+2

However, the modification is not placed in memory until the user types one of the three terminating characters - up arrow, backspace, or carriage return. The effect of each of these characters is given in the following table:

Terminating Character

Action

)

Returns carriage and modifies the contents of the open register if a modification has been typed. The register becomes closed. If a vertical bar was used to open the register, bits 2 through 5 indicate the core module that becomes "current".

←

Same action as carriage return except in addition the next sequential register in the current field is opened automatically (i.e., current field plus the address is typed followed by a slash tab, and the contents of the register). If no register is open when the backspace is typed, the next sequential register in the current field is still opened.

NOTE: If the current location is 7777, register 0 of the current field will be opened next.

↑

Same action as backspace except this character opens the preceding register of the current field instead of the following one.

NOTE: If the current location is 0, register 7777 of the current field will be opened next.

Once a particular register has been closed by use of either the carriage return, backspace, or up arrow, further modifications of that register is impossible until it is opened again.

4. Additional Interpretation of Register Contents - If, while a register is open, any one of the following characters is typed, the contents of that register will be reprinted in the indicated manner.

<u>Character</u>	<u>Interpretation</u>
=	types out quantity as a constant
→	types out quantity as an instruction
~	types out as if quantity is a concise code

To illustrate the use of these interpretation characters, consider the following examples:

```
reg+100/ →| lac abc →| = →| 202147 →| lac abc  
reg+101/ →| dac 6251 ~ →| ubr
```

where abc has the value 2147.

5. Examination and modification of a deferred register - Once an instruction has been typed out by ID, it is frequently desired to know the contents of the register addressed by the instruction. The control characters tab (→|), greater than (>), and special uses of slash (/) and vertical bar (|) provide this facility.

- a. After opening a register, the character tab (→|) may be typed to close that register and open the register in the current field addressed by its instruction. This causes the location counter, a register internal to ID which contains the address of the last register opened, to be changed. An example follows:

```
200/ →| lac abc →| →|  
abc/ →| 30 →| ←  
abc+1/→| 0 →|
```

Modifications may be made to a register while it is opened during this procedure. For example:

```
200/ →| lac abc →| lac abc+1 →|  
abc+1/→| 0 →| 5
```

- b. Like tab, the character > can be used to find out the contents of the register in the current field addressed by its instruction. Unlike tab, it just prints a tab and the contents (not the address followed by a slash). It opens, modifies, and closes registers in the same manner as tab. The current location counter is not changed. For example:

```
200/ →| lac abc →| ≥ 30 →| ←  
201/ →| dac bop
```

- c. The character / when used while a register is open closes the register without making any modifications to it and types out the contents of the register in the current field which was last typed by you or ID. The location counter is changed to the new register opened. For example:

```
200/ →| lac abc →| / →| 30 →| ←  
abc+1/ →| 0→|
```

or

```
200/ →| lac abc →| 100/ →| dac 1 →|
```

- d. Like /, the character | when used while a register is open closes the register without making any modifications to it and types out the contents of the 16-bit addresses register which was last typed by you or ID. The location counter and the current field are both changed according to this new register opened. For example:

```
1200/ →| 1+def →| | →| 2150 →| ←  
1 def+1/ →| 1567 →| (1=10000)
```

Notice that core 1 was made the current field and the location counter was changed to core 1 location def.

D. The Current Location Counter

The current location counter is a register internal to ID which contains the address of the last register opened in the current field. To re-open a register that has accidentally been closed or to refer to registers near the one presently opened, the current location character, point (.), is used. Typing an address followed by a register opening character such as slash or vertical bar sets the current location counter to that address. Backspace, up arrow (↑), and tab automatically set this register to the appropriate address; carriage return does not affect it. Since point (.) has the value of the current location, expressions such as `dap .+1` may be typed into ID (although they will not be typed out in this format).

E. Symbols and the Symbol Table

1. A symbol is a string of not more than six letters and numerals, containing at least one letter, and having a value associated with it. ID maintains a table of symbols and their values, and uses it to interpret symbolic words.
2. Initially, ID's symbol table contains 114 symbols, corresponding to PDP-1 instruction mnemonics, such as the operation mnemonics like lac, tyo, etc., the indirect bit i, the shift mnemonics 1s, 2s, etc.
3. There are five different ways of adding six character symbols to ID's symbol table.
 - a. A binary symbol tape may be prepared by POSSIBLE or MIDAS SYMBOL PUNCH and entered into ID by typing 1T. This causes the tape to be read and merges the symbols with ID's symbol table.
 - b. An alphanumeric or numeric tape may be prepared by POSSIBLE and entered into ID by typing T. This causes the tape to be read and merges the symbols with ID's symbol table.
 - c. Symbols may be defined directly by means of a close parenthesis as in the following example:

2475 sym) →|

The value 2475 is then associated with sym. Symbols may be redefined in this manner. (Even the initial PDP-1 mnemonics may be redefined, but there is rarely any reason to do so.) The redefinitions can be in terms of their old value:
If

abc=50

the command

abc+5 abc)

will make

abc=55

A symbol may be defined while a register is open by also using the close parenthesis. This would define the symbol to be the contents of the open register. For example,

1/ →| 720307 →| dpy)

defines dpy to be 720307.

- d. Symbols may be defined to be equal to the current location by typing the symbol followed by a comma. This does not affect the contents of the current location. For example, if the register last opened was 50:

50/ →| lac 774 →|

by typing

sym,

sym is defined as 50 but the register still contains lac 774.

sym/ →| lac 774 →|

- e. Symbols may also be defined to be equal to the 12-bit address part of the last expression typed by the user or ID by typing the symbol followed by an imply sign (▷). Thus:

500/ →| add 256 →| con▷

/ →| add con

Thus, con was defined to be 256.

4. Symbols may be destroyed by using the commands K or symK where sym is a symbol. The command K kills all the symbols in ID's table except the 114 PDP-1 instruction mnemonics. (If any of these were redefined, however, the original value is not restored.) The command symK kills only the particular symbol sym.
5. If a symbol which has not been defined is typed, ID types a capital U (undefined) and forgets the symbol.

F. Typing Instructions, Constants, and Location

1. Instructions, constants, and locations, which collectively may be referred to as words, may be typed by the user at any time using any combination of numbers and/or defined symbols separated by appropriate connectives such as plus and minus signs. In ID, a symbol is any combination of letters and numbers not longer than six characters, which contain at least one letter. (In most other versions of DDT, symbols can not be longer than three characters.)
2. The connectives used in forming words are listed in the following table along with their meanings.

<u>Connectives</u>	<u>Meanings</u>
■	adds value of next symbol or number to word.
+	adds value of next symbol or number to word.
-	subtracts value of next symbol or number to word.
^	ands value of next symbol or number onto word.
v	ors value of next symbol or number into word.

Thus,

<u>Typing</u>	<u>Yields</u>
add 10	400010
lac 2147	202147
lac i adr (where adr has previously been defined as 200)	210200
dpy-1	720007
clavcllvclfv _{space}	764207
law 144	700144

G. Evaluation of Words

1. Often it is desirable to be able to evaluate a word that is to be used in a program without actually affecting memory. This may be done at any time without opening a register by simply typing the word to be evaluated followed by the appropriate interpretation characters (see section C-4). When this is done, ID will automatically type out the appropriate interpretation of the word followed by a carriage return.

H. Notes on Symbolic Type-Out

A given register, containing only an octal number, can be interpreted symbolically in more than one way. Thus, ID may sometimes type out instructions you may not expect.

1. If several symbols are defined as having the same value, ID chooses to print out the last one defined. If clf 6 is typed into ID, it will be printed back as opr 6.
2. Symbols of four characters or more will only be printed out as the first three characters. Thus, if two symbols abc and abcd are defined as different octal values, ID will print both of them symbolically as abc. One may type the interpretive character to find which is used.
3. Expressions with negative terms will not type out as they were typed in; for example, if $ret=adr+5$, then ret-1 typed in will be typed out as $adr+4$. Similarly, ID recognizes the current location symbol (.) but never prints it out.

4. The symbols 1s, 2s, 3s, 9s are defined, but have been placed in a special part of the symbol table so as to be printed out only on shift and rotate instructions.
5. Operate group instructions and skip group instructions type out with inclusive or signs when necessary; for example, 762407 types out as clavcllvelf 7. Thus, if a register contains data which happens to be in this range, the resulting type-out may be in terms of these instructions.
6. Numbers beginning in 77_ _ _ type out as negative.

I. Control of Modes

1. Although it has been assumed so far that ID normally prints out the contents of registers as instructions with symbolic addresses and normally interprets constants as unsigned octal numbers, a provision has been made to alter this state of affairs with a considerable degree of flexibility.
2. There are several different register opening characters from which to choose according to the type-out mode desired.

Register Opening Characters

Meaning

/

Types out the contents of the preceding 12 bit address number as symbols or constants, according to the mode.

|

Types out the contents of the preceding 16 bit address number as symbols or constants, according to the mode.

[

Types out the contents of the preceding 12 bit address number as constants, but does not change the mode.

Register Opening
Characters

Meaning

]

Types out the contents of the preceding 12 bit address number as symbols, but does not change the mode.

(

Does not type out the contents of the preceding 12 bit address but puts ID into the type-in mode starting at that address.

3. Type-Out Mode For Instructions - By typing one of two commands to ID, the normal mode of printout of register contents may be controlled.
 - a. Symbolic type-out mode is the most often used and the one in which ID is initially. This mode is obtained by typing a capital S. The contents of registers will be printed out as symbols.
 - b. Constants type-out mode is obtained by typing a capital C. In this mode, the contents of the registers are typed out as numbers.

4. Type-In Mode is obtained by opening a register with an open parenthesis ((). In this mode, ID does not print out the contents of the register at all; it is a convenient mode for typing short programs or parts of programs. This mode is left by typing a carriage return; however, backspace, up arrow, and tab keep ID in type-in mode and open the appropriate register.

5. Type-out Mode for Address of Registers - By typing one of two commands to ID, the mode of printout of register addresses (as a result of tab, backspace, up arrow, etc.) may be set.
- a. Relative mode is the one in which ID is initially. By typing capital R the mode can be obtained again so that addresses will be typed out symbolically.
- ex. adr+10/ →| lac abc →| ←
adr+11/ →| dac x42 →|
- b. Octal mode causes the register addresses to be typed out as numbers. It is obtained by typing a capital O.
6. Constant print control - By typing one of the following two commands, the normal mode for the printout of constants may be controlled.

<u>Command</u>	<u>Resulting Action</u>
H	all constants will be printed out as octal numbers - <u>hoctal</u> mode.
U	all constants will be printed out as decimal numbers - <u>unhoctal</u> mode.

J. Input Radix Control

The commands H and U discussed in the preceding section control the radix used by ID to interpret all numbers typed out by the users. The character period (.) is used to force interpretation of input constant as decimal regardless of the current radix. If the input constant is not immediately followed by a period, it is interpreted as octal regardless of the current constant radix. The character single quote (') causes the last three characters typed in to be taken as their sqoze code value. This applies only to letters or numerals. The character double quote (") causes the first three characters typed in to be taken as their sqoze code value. This applies only to letters or numerals.

K. Special Registers

The capital letters in the following table indicate special consecutive registers, which are internal to ID. These registers control some of the main functions of ID; they may be referred from any field and are opened and modified in the same manner as a register in the current field.

<u>Capital Letter in their Location Order</u>	<u>Register Contents</u>
A	the stored accumulator of the program
I	the stored IO register of the program
X	the location of ID's execute register
G	the stored program counter of the program; the overflow flip-flop is stored in bit 0, the extend mode in bit 1.
F	the stored flags of the program
M	the mask for word searches
M+1	the lower limit for word searches, save and unsave fields, and special uses of yank, tape, verify tape, punch data blocks, and zero memory.
M+2	the upper limit for word searches, save and unsave fields, and special uses of yank tape, verify tape, punch data blocks, and zero memory.
B B+1 B+2 B+3	} breakpoint locations

The characters A, I, M, and B when preceded by a single argument deposit the argument in the corresponding register. For example, typing

17777A

deposits 17777 into ID's internal register A, containing the stored accumulator for the program.

The usage of the above control characters will be more fully explained in the sections to follow.

L. Assignment and Deassignment of IO Devices and Drum Fields

ID can assign or deassign IO devices and drum fields independent of the user's program. In this way, the user can be assured of obtaining essential equipment before starting his program running. The capital letter F when preceded by one or two arguments provides a convenient way to assign or deassign equipment directly from ID.* The table of possible request is found in the Assignment and Deassignment of In-Out Equipment and Drum Fields Memo (PDP-31). The mnemonic or concise code indicating the device requested is the argument immediately preceding the F command. Thus, typing αF where α is a mnemonic or concise code indicating the device requested, will request assignment of that device to the user's console. αF is equivalent to executing the following three instructions:

```
law flexo $\alpha$  or law $\alpha$ 
cli
arq
```

(NOTE: The arq is executed without reference to the special internal registers A and I of ID.)

In certain cases the IO must contain additional information about the device; thus the F command must have two arguments. Typing

$x\alpha F$

will put x into the IO and the concise code for the mnemonic of the device requested into the AC and then execute an arq. If the assignment or deassignment of

* The capital letter F when not preceded by an argument refers to ID's special internal register, F, containing the stored flags of the user's program.

fields is successful, then two carriage returns will occur. On the successful field assignments that return information in the AC, ID prints out the information in the right 6 bits and restores the AC to its contents before the request was made. For assignments and deassignments of in-out devices (not fields) two carriage returns will be returned only on successful assignments. [On unsuccessful assignments, only one carriage return is given.] An assignment will be successful if the field(s) or device requested is not already assigned or if the assignment is already in effect.

M. Breakpoints

One of the most powerful features of ID is the ability to insert breakpoints in programs. In testing a large program, it is frequently convenient to use breakpoints to interrupt the computation so that partial results may be examined or the state of the program determined. Breakpoints may be set up at a location in the user's program by two methods:

1. Typing

adrB

causes ID to set up a breakpoint in the current field at location adr. Only one breakpoint can be inserted at a time by this method; the address preceding the B will be deposited into the special register B.

2. Four special registers, B, B+1, B+2, and B+3, can be used to contain the addresses of breakpoints. No break location is indicated by an overbar ($\bar{\quad}$); initially all four registers contain overbars. For example:

B+1/ \rightarrow | $\bar{\quad}$ \rightarrow | adr

This puts a breakpoint at location adr in the user's program. If the user transfers control to his program,

and the instruction in register adr is reached, computation will cease and control will be returned to ID, which will type out the register location, a close parenthesis, tab, and the original contents of the register. At this point, the user may examine the accumulator, IO, and/or any other register and make modifications as he pleases. A breakpoint remains in the location specified until it is removed by clearing the breakpoint register containing the address. All breakpoints may be cleared by typing B⁻. If the user wants to clear only one breakpoint, he puts an overbar or a minus zero in the breakpoint register containing the break address to be cleared.

CAUTION: The location selected as breakpoints must not be registers whose contents are modified by the program under test, since ID transplants their contents and substitutes specific transfer commands.

N. Go (G), Proceed (P), and Execute (X)

1. The instruction adrG, where adr is an address in the user's program, is used to start the user's program running at location adr.
2. If a breakpoint trap occurs, control is transferred to ID. To continue operation of the user's program from the point at which the break occurred, the command P is used. Even if the last breakpoint encountered has been deleted or moved, P still proceeds from the point where the break actually occurred.
3. nP, where n is a positive numeral, will cause ID to proceed from a breakpoint trap, and go past the breakpoint n times before breaking again. This multiple proceed command works only for the breakpoint whose address is in register B.
4. Single instructions may be executed directly by ID; control need not be returned to the user's program. There are two possible ways to execute single instructions in ID:

a. Typing

bX

causes the instruction b to be placed in the address specified by the contents of the execute internal register X and then to be executed.

b. Typing

a<bX

causes the instruction b to be placed in address a and then to be executed. The internal register X does not change.

Normally there are two carriage returns after X; if the PC is incremented by two (that is, the instruction skips), X will return the carriage a third time. If the return PC is not the same as the original PC incremented by one or two, control is transferred to the location specified by the PC. Otherwise control is returned to ID.

0. Word Searches

A valuable feature of ID is its search facility. Three kinds of searches can be made; these types are controlled by the commands N, W, and E, and they all use the special internal registers M, M+1, and M+2.

1. The three types of searches and their respective commands are:
 - a. wordW - The word search causes ID to search the current field for and print out all the registers, between the limits in M+1 and M+2 inclusively, containing the given word.
 - b. wordN - The non-word search causes ID to search the current field for and print out all the registers, between the limits in M+1 and M+2 inclusively, not containing the given word. This is most frequently used in ON, the search for non-zero memory.
 - c. adrE - The effective word search causes ID to search the current field for and print out all the registers, between the limits in M+1 and M+2 inclusively, effectively addressing adr. If the user is in extend mode, (bit 1 of the PC on), indirect addressing chains for effective address searches will be carried to a depth of 1; otherwise they will be carried to a depth of 10, at which point ID will give up.* An E search will never print out skp, sft, law, lot, 74, or opr instructions. This type of word search is valuable for locating incorrect instructions which are modifying the program. If a jda instruction is suspected, try jda adrN in addition to adrE.

* An E-search with greater depth than 10 octal might take a long time and an E-search with no restriction on depth might get caught in an infinite chain like:

```
adr,          lac i abc
abc,          jmp i adr
```

2. The special internal registers for word searches are M, M+1, and M+2; the use of these registers is explained in the following table.

<u>Register</u>	<u>Contents</u>
M	The mask register contains the value of the mask used in word searches. During word searches, only the bits masked 1 in register M are compared. Initially M contains -0; thus all bits are compared unless the register is modified.
M+1	The lower limit for the word search is stored in the M+1 register. Initially, M+1 contains 0; thus the search will begin at 0 unless modified.
M+2	The upper limit for the word search is stored in the M+2 register. Initially, M+2 contains 7777; thus the search will end at 7777 unless modified.

3. Special commands may be used to modify the contents of the special internal registers M, M+1, and M+2.

Typing

M

initializes the contents to -0 in M, 0 in M+1, and 7777 in M+2.

fa<laM

puts fa and la in M+1 and M+2 respectively. M remains unchanged. To change M, type

aM

where a is the mask desired for M.

4. There are two ways to print a block of registers:
 - a. Set the mask to zero and set up M+1 and M+2 to enclose the area to be printed. Then search for any word.
 - b. If irrelevant parts of memory happen to contain zero, merely do a N- search for zero.

P. Zero

Often it is valuable to zero all or parts of a field so that irrelevant parts of the field will contain zero.

The following commands may be used:

<u>Command</u>	<u>Meaning</u>
Z	zero all of the current field
fa<laZ	where fa and la are 12-bit addresses limits for the zero command. The registers of the current field between fa and la inclusively are zeroed by this command.
xZ	where x is the field number for the zero command. The field specified is zeroed between locations in M+1 and M+2 inclusively. The current field is not changed.

Q. Yank

In the Preparation Section of this memo (part A), the user was instructed to use the command "Y" to read into the current field a binary tape prepared by POSSIBLE or MIDAS. For convenience, other variation of this command may be used. They are:

Command

Meaning

Y

Read a tape in POSSIBLE or MIDAS binary block format into the current field between the locations specified by M+1 and M+2 inclusively. The core modules specified in the data block origins will be ignored. If a checksum is encountered, the process will stop. It is then possible to move the tape back one block, restart the reader, and type "c" to continue reading, if desired.

xY

x is the field number into which a tape in POSSIBLE or MIDAS binary block format is read. Otherwise, the command is the same as Y alone. The limits of the yank are in M+1 and M+2 as above. The core modules specified in the data block origins will be ignored. If a checksum is encountered, the process will stop. It is then possible to move the tape back one block, restart the reader, and type "c" to continue yanking, if desired.

fa<laY

where fa and la are 16-bit address limits for the yank command. The data block will be checked against core field specified in the block origin. Only words with extended addresses from fa to la inclusively will be yank in. The process will stop on encountering a checksum. To continue, move the tape back one block, restart the reader, and type "c" to continue yanking.

R. Verify

Another feature of ID is the ability to verify the program currently in core or on a drum field with the original binary tape. The capital letter V is used as the command in the following ways:

Command

Meaning

V

Read a binary tape in POSSIBLE or MIDAS binary block format; the core modules specified in the data block origins will be ignored. The words read in are compared against the current fields words between locations specified by M+1 and M+2 inclusively. No change is made to memory; any discrepancies are typed out as:

location/ memory tape

If a checksum error is encountered, the process will stop. It is then possible to move the tape back one block, restart the reader, and type "c" to continue reading, if desired.

xV

x is the field number whose contents is to be compared against the tape. The field may be a core field or drum field. Otherwise, the command is exactly the same as V alone. The limits of the verify are in M+1 and M+2 as above. No change is made to memory and any discrepancies are typed out as:

location/ memory tape

The program will stop on encountering a checksum. To continue, move the tape back one block, restart the reader, and type "c" to continue reading and verifying.

Command

Meaning

fa<laV

where fa and la are 16 bit address limits for the verify command. The data blocks will be checked against core field specified in the block origin. Only words with extended addresses from fa to la inclusively will be checked. No change is made to memory and any discrepancies are typed out as:

extended location memory tape

The process will stop on encountering a checksum. To continue, move the tape back one block, restart the reader and type "c" to continue verifying.

S. Save and Unsave Drum Fields

Another valuable feature of ID is the ability to save an image of a program on another drum field, so that it may be restored at some future time. The capital letters S and U, when preceded by additional information, are command to save and unsave drum fields.* The special internal registers M+1 and M+2 indicate the limits of the transfer for the current field. The two basic commands and their meanings are:

<u>Command</u>	<u>Meaning</u>
fS	Save on field "f" - an image of the current field between the limits in M+1 and M+2 is written onto drum field f between the limits also in M+1 and M+2. This operation does not affect the contents of the current field. Field "f" must be assigned to your consoles; it must be a number from 1 to 17 ₈ when referencing a pseudo field, or from 41 ₈ to 57 ₈ when referencing an absolute field.
fU	Unsave field "f" - the contents of the current field between the limits in M+1 and M+2 are replaced by the contents of drum field f between the limits in M+1 and M+2. The contents of drum field f are not affected by this operation. Field "f" must either be an absolute system field or a field assigned to your console; thus it must either be a number from 1 to 17 ₈ when referencing a pseudo field assigned to your console, a number from 41 ₈ to 57 ₈ when referencing an absolute field assigned to your console, or a number from 61 ₈ to 66 ₈ when referencing an absolute system field.

* The capital letters S and U when not preceded by a character mean symbolic and unhoctal. (See section I-3 and 6.)

Two other commands to unsave and save drum fields are available for swapping information to a different location on the drum field and the current field. These commands are:

Command

Meaning

x<fS

Add "x" to the origin of the area on field f - an image of the current field between the limits in M+1 and M+2 is written onto drum field f between the limits "x" plus the contents of M+1 and "x" plus the contents of M+2. Thus the limits in M+1 and M+2 apply only to the current field, not field "f". Field "f" must be assigned to your console; it must be a number from 1 to 17₈ when referencing a pseudo field or from 41₈ to 57₈ when referencing an absolute field.

x<fU

Add "x" to the origin of the area unsaved from field "f" - the contents of the current field between the limits in M+1 and M+2 are replaced by the contents of drum field f between the limits "x" plus the contents of M+1 and "x" plus the contents of M+2. Thus, the limits in M+1 and M+2 apply only to the current field, not field "f". Field "f" must either be an absolute system field or a field assigned to your console; thus, it must either be a number from 1 to 17₈ when referencing a pseudo field assigned to your console, a number from 41₈ to 57₈ when referencing an absolute field assigned to your console, or a number from 61₈ to 66₈ when referencing an absolute system field.

An example of using the latter commands appears below:

100<200M

20<5S

moves locations 100 - 200 inclusive from the current field to locations 120 - 220 of field 5. To restore this program material at a later time, the user would type:

100<200M

20<5U

and thus move locations 120 - 220 of field 5 to 100 - 200 of the current field.

T. Hoarding and Reading Symbols

Another feature of ID is the ability to hoard and read symbols, so that the symbols may be stored and restored with the associated program. The capital letters H and R, when preceded by additional information, are commands to hoard and read symbols.* The two basic commands and their meanings are:

Command

Meaning

fH

Hoard ID's symbol table on field f - saves all of the user's symbols (except initial symbols, even if redefined) on the part between n and 7777 inclusive. The number n is printed out and becomes the new memory bound for field y. (N is also in location 7777.) This feature is intended to be used in association with "S" to save a program on lower portion of a field and its associated symbols on the upper portion of the same field. The symbols are not changed or killed in any way by "H". Any argument acceptable to "S" as a field number is acceptable to "H".

fR

Read the symbols stored on field f into ID's symbol table - reads all of the user's symbols previously stored on field f by the command "H" and bodily appends it to ID's initial symbol table. Previous symbols in ID's symbol table are killed (except initial symbols). If what it finds on that field is not a symbol table, it responds with "?", and ID's symbol table is not changed. This feature is intended to be used with "U" to unsave a program and its associated symbols for further reference. Note that the "R" process is different from "T" in that in case of "R", current symbols are first killed, whereas in the case of "T" new symbols read are merged with current ones. Any argument previously used by "H" as a field number can be used for "R".

Two other commands to hoard and read symbols are available for swapping the symbols to and from a specified location. These are:

Command

Meaning

x<fH

Hoard symbols on field f below location x. The number n is printed out; the table of user's symbols is between n and x-1 inclusively. (N is also in location x-1.) x may be any symbolic or numeric location and any argument acceptable to "S" as a field number may be used for f in this command. The symbols are not changed or killed in any way by this command.

x<fR

Reads symbols from field f below location x previously stored by x<fH and appends them to ID's initial symbol table. Previously symbols in ID's symbol table are killed (except initial symbols). If what it finds on that field is not a symbol table, it responds with a "?", and ID's symbol table is not changed. Any arguments previously used in the "x<fH" command can be used for "x<fR".

* The capital letters H and R when not preceded by a character mean hoctal and relative. (See section I-5 and 6.)

U. Punching Programs

When final corrections have been made in the user's program, the user may punch it out in its modified form. The four punching commands are L, D, center dot, and J.

1. L causes ID to listen for title. Letters typed after this command will be punched in readable form on tape. The title punch is terminated by carriage return, tab, or backspace. The result of these terminating characters is given in the following table:

<u>Terminating Character</u>	<u>Result</u>
)	Punches the standard input routine and sets ID to punch the usual checksummed data blocks.
->	Sets ID to punch the usual checksummed data blocks, but no input routine.
←	Sets ID to punch read-in mode tapes.

2. The capital letter D is used to punch data blocks from the current field. A variety of formats are available to the user for his convenience.
 - a. fa<laD, where fa and la are any symbolic or numeric expressions, punches the current field from fa to la inclusive. If the current field is a drum field, the origins of the data blocks will be in core 0. If the current field is a core field, the origins will be in the current core.
 - b. D alone is equivalent to 0<7777D. It punches the entire current field. If the current field is a drum field, the origins of the data blocks will be in core 0. If the current field is a core field, the origins will be in the current core.
 - c. xD, where x is a core number 0 to 7, punches the current field between the limits in M+1 and M+2. The data block origins will be in core x.

3. aJ, where a is any symbolic or numeric expression, causes ID to punch a start (jump) block to the address specified to denote end of binary tape. The address is typed immediately preceding the J.

4. If a register is open, center dot (•) will close the register and punch its contents as a one-word data block. This is convenient if the tape needs only a few modifications, known in advance.

V. Error Indications and Corrections

1. ID has several error alarms associated with its use. These are typed out by ID and have the following general meanings:

cksm	A sum check error occurred in reading a binary program or symbol tape. By moving the tape back one block and typing "c" the tape will continue to be read.
d.e.	Drum swap was not successful. Error may be caused by trying to write on locked field, or a timing error in drum.
no L?	This indicates that user tried to punch out before obtaining a title on the tape.
busy	This indicates that the reader or punch is busy and the user must wait until available.
U	This indicates that the immediately preceding word contains an undefined symbol. ID will act as if nothing had been typed. Thus, for example, typing an undefined symbol in a word into an open register will result in "U", but typing a carriage return will close the register with its previous contents rather than zero.
?	Error has been made in the command to ID. ID can't do or doesn't understand the request typed in.

2. When a user's program executes an illegal instruction, ID is brought back into control and the address of the illegal instruction is typed and followed by >> and a tab. Then, the contents of that register are typed out. Below is a list of various types of illegal instructions:
 - a. hlt instruction
 - b. instruction with an illegal operation code.
 - c. instruction which directly or indirectly addresses a location above the memory bound.
 - d. a reader or punch instruction when no assignment has been obtained for the program.
 - e. arq instruction with invalid code or parameters.
 - f. a dcc drum instruction addressing an unassigned field or locations in core above the memory bound.
 - g. a bpt instruction at a location to which a breakpoint was not assigned by the user through ID.

3. When the user of ID realizes that he has made a typing error, he may delete all that he has typed since the last carriage return or tabulation by typing a multiplication sign (x). For example:

```
loc/ →| add a →| abcx →| add abc
     _/ | add abc.
```


APPENDIX I
SUMMARY OF CONTROL CHARACTERS

A.	accumulator storage (19)*
B,B+1,B+2,B+3	registers containing breakpoint locations (22)
C	set word print mode to constants (17)
D	punch data blocks (36)
E	effective address search (25)
F	without argument: storage for program flags (19) with one or two arguments: execute an arq (21)
G	without argument: storage for program counter (19) with one argument: start program running, go to (24)
H	without argument: set constant printout mode to (18) hoctal (octal) with one or two arguments: hoard symbols onto field(34)
I	i-o storage (19)
J	punch start (jump) block (37)
K	kill defined symbols (3)
L	listen for title punch (36)
M	mask register (26)
M+1	lower limit for word search (26)
M+2	upper limit for word search (26)
N	not-word search (25)
O	set location print mode to octal (18)
P	proceed (24)
Q	last quantity (7)
R	without argument: set location print mode to (18) relative with one or two arguments: read symbol table (34) from field.
S	without argument: set word print mode to symbolic (17) with one or two arguments: save memory on field (31)
T	read symbol table (T, 1T, 2T) (12)

* The numbers in parentheses indicate the page number where the character can be found.

U without argument: set constant printout mode
to unhoctal (decimal) (18)
with one or two arguments: unsave field into (31)
current field

V verify tape (29)

W word search (25)

X execute as instruction (24)

Y read binary tape (28)

Z zero memory (27)

0-7 octal numerals and/or symbol constituents (14)

8,9,a-z symbol constituents (12)

" take as concise code (18)

~ print as concise code (9)

D define symbol as address typed (13)

V inclusive or (14)

^ and (14)

↑ modify and open previous register (8)

→ print as instruction (9)

(open register in type-in mode (17)

) define symbol (12)

[examine register as octal constant (16)

] examine register as instruction (17)

- minus (14)

+ plus (14)

⊠ (space) plus (14)

,

= print as octal (9)

.

current location; if preceded by number take
constant as decimal integer (18)

x delete typed input (39)

/ examine 12-bit address register (10)

tab	modify and open addressed register; also alters sequence of location (9)
bk sp	modify and open next register (8)
car ret	modify and close register (8)
uc, lc	set case
	examine 16-bit address register (10)
>	modify and open addressed register (10)
'	take preceding constant as decimal integer (18)
• (center dot)	punch opened register as one word block (37)

APPENDIX II
ID SYMBOL TABLE

BASIC INSTRUCTIONS

add	400000
and	020000
cal	160000
dac	240000
dap	260000
dio	320000
dip	300000
dis	560000
→div	560000
dzm	340000
idx	440000
lor	040000
lot	720000
isp	460000
jda	170000
jmp	600000
jsp	620000
lac	200000
law	700000
lio	220000
→mul	540000
mus	540000
opr	760000
sad	500000
sas	520000
→sft	660000
skp	640000
sub	420000
xct	100000
xor	060000

OPERATE GROUP

cla	760200
clc	761200
clf	760000
cli	764000
cma	761000
hit	760400
→lai	760040
lap	760300
lat	762200
→lia	760020
nop	760000
opr	760000
stf	760010
→swp	760060
xx	760400

SKIP GROUP

→clo	651600
skp	640000
sma	640400
sni	644000
spa	640200
spi	642000
→spq	650500
sza	640100
szf	640000
→szm	640500
szp	641000
szs	640000

IN-OUT TRANSFER GROUP

cbs	720056
cks	720033
→dba	720061
→dec	720062
→dia	720060
dpy	730007
→dra	720063
eem	724074
esm	720055
ioh	730000
lot	720000
lem	720074
lsm	720054
ppa	730005
ppb	730006
rpa	730001
rpb	730002
rrb	720030
tyi	720004
tyo	730003

SHIFT/ROTATE GROUP

ral	661000
rar	671000
rcl	663000
rcr	673000
rll	662000
rll	672000
sal	665000
sar	675000
scl	667000
scr	677000
→sft	660000
sll	666000
sir	676000

MISCELLANEOUS

→clo	651600
i	010000
is	1
2s	3
3s	7
4s	17
5s	37
6s	77
7s	177
8s	377
9s	777

TIME SHARING INSTRUCTIONS

sps	723077
sd1	723477
lsb	720052
wat	722477
arq	722277
bpt	722177
dsm	722377
ckn	720027
rbt	720217
cac	720053
asc	720051
dsc	720050
nmn	725377
nmf	725477
sbr	722577
lea	724677
lei	724577
rer	724777

ID 16 aug 65 part 0

2200/

bpt=iot 2177
arq=iot 2277
dsm=iot 2377
jdp=140000
adm=360000

/low end initial symbol table
low=-2400+2033
/top end of symbol table
tst=-2
/end of symbol table, pointer t low end
est=-1
/symbol limit, max number. est-sl-sl must be \geq
/lowest location used by symbol table is est-sl-sl-2.
sl=1000
/number of break points
nbp=4
/number of internal registers
nir=nbp+6
/number of drum fields
ndf=26

/fixed loc in exec routine
xlc=16

2200, jmp ent /begin at the beginning, the king said....

```

define feed a
law i a
jda fee
        terminate
define dispatch low,upp
[upp-uc 44]x1000 low-uc 44
termin
define letter a,b
disp [[a+uc-44]],b
        terminate
define bprint
b=1
repeat nbp-1,746254
        terminate
        /good for up to nbp=10
        7200+b      b=b+1

```

```

dfp,      add      /drum field of program (real)
dpf,      add      /used for holding dfp

```

```

aa,      746100      0      /A
          747100      0      /I
          746600      0      /F
          744400      0      /M
          744454      007201  /M+1
          744454      007202  /M+2
          746200      0      /B
          bprint      /B+1, etc.

```

```

l=0
r=1

```

```

start

```

ID part 1 1/26/65.

/this is the beginning of the real program

```
ent,      lac ids   /were we in ID?
          sza
          jmp en1   /yes
          idx ids   /we are now
          jsp ssw   /set state word pointers
          eem
          lac i acp
          dac ac
          idx acp
          lac i acp
          dac io
          idx acp
          lac i acp
          dac pc
          idx acp
          lac i acp
          dac fg
```

/now get the breakpoint package

```
          law bp
          dap . 2
          lac i bpp
          dac bp
          idx bpp
          idx .-2
sas (dac bp nbp
          jmp .-5
```

/get kdm pointer

```
          law xlc
          dap xmw
          lac i xmw
          add xc0   /kdf-aw1 = 4
          dap xmw
          lac i xmw
          lem
          dap kdm   /pointer into exec core
```

/get why

```
          jsp gy
          jmp en2
```

```

en1,      jsp gy
          sad (4      /call button
          jmp zrt
          lac dpf
          dac dfp      /restore drum field
          jmp err

```

/dispatch

```

en2,      and (77
          sas why      /check validity
          jmp .         /Leo, you goofed.
          add red
          dac bkf      /setup flag
          dap . 1
          jmp .         /go, man, go

```

/dispatch table

```

red,      . 1          /constant
          jmp zrt      /'ddt' typed from mystic
          jmp .         /arq--leo goofed again
          jmp cl       /dsm issued
          jmp ii       /illegal op
          jmp cll      /call button
          jmp lbp      /bpt given
          jmp rx       /return from X
          jmp .         /leo screwed up

```

/subroutine to get why, the reason ID was entered
/leaves reason in why and ac - io preserved

```

gy,       dap gyx
          eem
          law xlc
          dap xmw
          lac i xmw
          add xc1
          add xc2
          sub xc3
          dap xmw
          lac i xmw
          dap xmw
          lac i xmw
          lem
          dac why
gyx,      jmp .

```


/random constants

```
acp,      050000  /pointer to ac in exec core
bpp,      050000  /pointer to bpt package in exec core
xmw,      050000  /gets into exec core.
xc0,      000004  /04=kfd-aw1 (executive symbols)
xc1,      000017  /15=sp2-aw1 (executive symbols)
xc2,      000015  /13=brk-sp2      "
xc3,      000003  /brk-why        "
why,      000000  /reason ddt was activated
bp,       000000  /break-
           000000  /point
cn1,      000000  /pack-
           000000  /age.
xe2,      000000  /temp storage for pc
nms,      000000  /switch for G P or X
ids,      000000  /ID switch(non 0-ID running)
opc,      000000  /storage for pc during X
kdm,      050000  /pointer to kdm word in exec
bpa,      000000  /temp for pointer tracing
xe3,      000000  /temp for instruction in bpt routine
p11,      000000  /storage for proceed routine
mpc,      lac bk1 nbp  /random number for bpi-bpo
mb,       006770  /memory bound
```

/subroutine to set up pointers to state word and bpt pack
/sets pointers acp and bpp
/jsp ssw to enter ac clobbered io preserved

```
ssw,      dap ssx
           law xlc
           dap xmw
           eem
           lac i xmw
           add xc1  /sp2-aw1 = 15
           dap xmw
           add xc2  /brk-sp2 = 13
           dac ssx 1
           lac i xmw
           dap xmw
           lac i xmw
           dap acp  /pointer to ac
           lac ssx 1
           dap xmw
           lac i xmw
           dap xmw
           lac i xmw
           dap bpp  /pointer to bpt package
           lem
ssx,      jmp .
           0      /temporary storage
```

cl, clc /clear the proceed indicator, dsm entry
 dac bk0
 jmp xe 3

cl1, law 2 /call button entry
 sas nms
 jmp cl1
 clc /call while X in progress
 dac bk0 /cant proceed
 law 7427
 jda tys
 law 7255
 jmp ixex 3 /pc=xe2 so mp3 works right

cl1, law 7777 /normal call
 and pc
 jmp 3bp

ixex, law 7427 /illegal execute
 jda tys
 law 1010
ixex+3, jda tys
 jsp lct
 jsp mp3 /setup
 jsp fet /and get the offending word
 jda lwt
 jmp xex+1

rx, lac pc /X return
 sub xex2 /did X
 sas (1 /skip?
xex, jsp lcc /yes
 lio opc /no skip
 dio bk0 /restore old pc for proceed after X
 jsp bpo
 jmp zrt

/drum read error routine

dre, lac (723564
 jda tys
 lac (6534
 jda tys
 jmp lis

```

1bp,      dzm bkf      /breakpoint
          law 2
          sad nms
          jmp ix     /bpt while X in progress.
          law 7777
          and bp 3   /true loc of bpt
          dac t
          law bk1    /setup
          dap 2bp
2bp,      lac .          /check for assigned breakpoint
          sad t
          jmp 3bp
          idx 2bp
          sas (lac bk1 nbp
          jmp 2bp

ii,       law 2        /illegal instruction
          sad nms      /were we X ing
          jmp ix     /yes-illegal execute
          lio (741010
          law 7777    /check
          and pc
          dap t2      /for
          law bk1
          dap . 1     /illegal
          lac bk1
          sad t2      /instruction
          jmp 3bp 1
          idx .-3     /under
          sas (lac bk1 nbp
          jmp .-5     /breakpoint.
          dio bkf
          lac t2      /to set up bk0 after bp3
          jmp 3bp 1

3bp,      lio (55
          dio t3
          dac bk0     /set up for proceed from bpt or illegal
          jsp bpo
          lac pc
          dzm mod
          dzm lcf
          jda pad
          lac t3
          jda tys
          lac pad
          jmp ta5

```

/G go

```
bg0,      jda  chk
          dzm  nms      /new mode switch
          spi
          jmp  err
          dap  xe2      /setup address
          jsp  mp3      /setup fetch and counter for bpt
          jsp  bpi      /put in breakpoints
          jsp  xxx      /for sequence break(same xxx as exec)
p1,       jsp  lcc
          jmp  p12
```

/X execute from location n

```
xe0,      spi
          jmp  err
          dap  xe2
          law  2
          dac  nms
          lac  bk0
          dac  opc      /save old pc til after X
          jsp  mp3
          jsp  fet
          dac  xe3
          jmp  p1
```

/subroutine mp3

```
mp3,      dap  . 7
          dzm  tsf      /fetch from drum
          lac  xe2
          add  c4
          dac  tas      /at loc=c(xe2)
          lac  c4
          dac  cn1      /proceed counter= 0
          jmp  .
```

/P proceed and multiple proceed logic

```
pr0,      spi
          law  0      /single proceed
          cma
          add  c4
          dac  cn1
          lac  bk0      /loc to proceed from
          sad  (-0
          jmp  err      /cant proceed
          dap  xe2
          dzm  nms
          law  bk1
          dap  p8
          dap  . 1
p2,       lac  .
          sad  bk0      /is this proceed through bpt?
```

```
jmp p3  
idx p2  
sas (lac bk1 nbp
```

```
jmp p2
```

```

lac bkf /did we stop by xct ing a bpt?
sza i
p8, jmp p1-2 /no
lac bk1 /yes
sad bp 3 /is there still one there?
jmp p9 /yes
idx p8
sas (lac bk1 nbp
jmp p8
jmp p1-2 /no longer there

p9, lac p8
jmp p3 1

p3, lac p2
dap p11
jsp bpi
jsp lcc
lac p11
dap . 1
lac .
add c4 /after putting in bpt s
dap bp 3
dac tas /put back old instruction
dzm tsf
lac p11 /so that it will be executed
add (nbp
dap . 1
lac . /this is the instruction from under bpt
jda dep /put it back
law 1
dac nms /setup proceed.
jmp p1-1

```

/xxx routine and associated trash

```

xxx, dap xx1
eem
lac fg
and (23000
sad (21000
law 1
lio i kdm
ril 1
rcr 1
dio i kdm
lem
xx1, jmp .

```

/set the state of the user and go

```
p12,      jsp ssw
          laaaaaaaccccc   bk1      lac bk1
          dac bp
          lac bk1 nbp
          dac bp 1 /set up bpt package
          eem
          lac ac
          dac i acp
          idx acp
          lac io
          dac i acp
          idx acp
          lac xe2
          dac i acp
          idx acp
          lac fg
          dac i acp
```

/now the bpt stuf

```
law bp
dap . 1
lac bp
dac i bpp
idx bpp
idx .-3
sas (lac bp 4
jmp .-5
dzm ids /note that we are leaving ID
lac nms
dsm /pass the buck to exec.
```

```

bp1,    lio (-0          /breakpoints insert into user's field
        jmp .+2
bp0,    cli          /break points take out
        dio dff
        dap bp6
        lac bp6+2
        dac bp2+1
        law bp5
        dap dpx
        spi i
        idx dpx
bp1,    law bk1
        dap bp3
bp2,    add (nbp
        dap bp4          /gets changed
bp3,    lac .          /used by cbp
        spa
        jmp bp5+1
bp4,    dap tas
        lac .
        dac dep
        cla
        jmp dp0
bp5,    dac i bp4
        idx bp3
        sas mpc
        jmp bp2
        cla
        sad dff
bp6,    jmp .
bp6+2,  law (bpt
        dap bp4          /used as a constant above
        dzm dff
        idx dpx
        lac .
        dac bp2+1
        jmp bp1

cbp,    dap cbx          /clear all breakpoints
        law bk1
        dap bp3
        clc
        dac i bp3
        idx bp3
        sas mpc
        jmp .-4
cbx,    jmp .

```


fet,	lio (-0 dio dff jmp dep+2	/fetch
dep,	. dzm dff dap dpx lac tas sma jmp dpx-1	/deposit
dep+6,	lac tsf sza jmp dp2	/re-entry from word search
dp0,	lio dff spi i lac dfp dip t1	/re-entry from bpi, bpo
dp1,	law 7777 and tas jda chk dap t1 lio t1 dia law 1 lio dff spi add dfp swp law dep dcc jmp dre lac dep jmp .	
dpx,		
chk,	0 dap chx law 7777 and chk sas chk jmp ta0 sub mb sma jmp ta0 add mb jmp .	/check against F
chx,		

```

dp2,    lio dff                /fetch, deposit internal registers
        law 7777
        and tas
        sub (ac
        spa
        jmp ta0
        sub (4
        sma
        jmp dp4
dp3,    lac i tas
        spi i
        lac dep
        dac i tas
        jmp i dpx
dp4,    sub (2
        sma
        jmp dp6
        spi
        jmp dp3
        add one
        sma
        jmp dp5
        lac dep                /M+1
        jda chk
        jmp dp3
dp5,    lac dep                /M+2
        jda chk
        sub ll
        spa
        jmp ta0
        jmp dp3
dp6,    sub (nbp
        sma
        jmp ta0
        spi                    /B thru B+nbp-1
        jmp dp9
        lac dep
        sad (-0
        jmp dp3
        jda chk
        dac chk
        law bk1
        dap dp7
dp7,    lac .                    /check whether already assigned
        sad chk
        jmp dp8
        idx dp7
        sas mpc
        jmp dp7
        jmp dp3
dp8,    clc
        dac i dp7
        jmp dp7+3

```

dp9, lac i tas
 sma
 jmp i dpx
 law 56
 jda tys
 clc
 dac lwt
 jmp pn2

pv, dap pvx /punch, verify swap routines

 clc
 dac dff
 lac pvf
 sza i
 jmp pv1
 spa
 jmp pv2
 law 7777
 and fa
 sub lo
 spa
 jmp pv1
 dap .+5
 sub wc
 add one
 szm
 jmp pv1

 / from pwd

hi, lac . /used by searches

pvx, jmp .
pv1, lac fa
 dac pvf
 dap lo
 jsp zd
 lac 0
 jmp pvx
pv2, lac dep
 jmp pvx

```

zd,      dap zdx      /zero drum, used for searches also
        law 1
        add wrd
        sub lo
        sub est
        sma
        cla
        add est
        dap esp
        dap wc
        lac dff
        sza
        jmp zd2
        dap .+1
        dzm .
        idx .-1
        sas esp
        jmp .-3

zd1,     lac lo      /re-entry point
        add dfp
        swp+cli-opr
        dia
        lio wc
zda,     dcc
        jmp dre
zdx,     jmp .
zd3,     law i 1     /re-entry point
        add lo
        add wc
        sub wrd
        sza i
        jmp zrt
        lac wc
        adm lo
        law 1
        add wrd
        sub lo
        sub wc
        sma
        cla
        adm wc
        lio dff
        spi i
        jmp zd1

zd2,     lio lo
        dia
        lac wc
        add dfp
        swp+cli-opr
        jmp zda

```

```

sav,      nop          /or sas mst+3 for field 23 for radio astronomy.
          jda ckf
          dac t1
          lia
          lac dfp
          dac t
          jmp su1
uns,      jda ckf      /unsave
          dzm t
          dip t
          lio t
          lac dfp
          dac t1
su1,      dzm lo
          law i 1
          add mb
          dac wrd
          clc
          dac dff
          lac t
          dac dfp
          jsp zd
su2,      lac t1
          dac dfp
          law su3
          dap zdx
          jmp zd1
su3,      lac t
          dac dfp
          law su2
          dap zdx
          jmp zd3

spf,      add (20
          cli+swp-opr
          rcr 6s
          add c4
          jmp su1-5

ckf,      0
          dap cfx
          lac ckf
          and (37
          sub (ndf
          szm
          jmp err      /field number too big
          lac ckf
          rar 6s
cfx,      jmp .

```

ws,

/W,N,E

spi
jmp err
dzm tsf
dap ws2
lac ll
sub ul
szm
jmp err
jsp lcc
clc
dac dff
lac wrd
dac chi
lac ll

dac lo
lac ul
dac wrd
jsp zd

ws0, lac wc
dap hi
cla
dap ws4+1

ws4, dzm sym
lac .
dac t2

ws2, jmp . /ea1 or ws1

ea1, and (770000
sad (jda
jmp .+4
and ci
sza
jmp ea2
law 7777
and t2

ws1, xor chi
can, and msk /used as and
wea, xx /sza or sza i
jmp ws3

law 7777
and ws4+1
add lo
dap loc
dzm lcf
jda pad
law 2136
jda tys
lac i ws4+1
jda lwt
jsp lcc

ws3, idx ws4+1
sas hi
jmp ws4
jmp zd3

ea2, idx sym
sad c10
jmp ws3
law 7777
and t2
sub mb
sma
jmp ws3
add mb
dac tas
sub lo
spa
jmp dep+6
dap .+5
sub wc
add one
szm
jmp dep+6
lac .
jmp ws4+2

```

ver,      jsp ar          /verify
          jsp lcc

vf1,      lac t
          dap fa
          lac fa
          sub chk      /10
          sub (dio
          spa
          jmp vf2
          add chk      /10
          sub wrd
          szm
          jmp vf2
          jsp pv
          dac chi
          lac i la
          sad chi
          jmp vf2

vf3,      lac fa
          dap loc
          dzm lcf
          jda pad
          law 2136
          jda tys
          lac chi
          jda lwt
          jsp lct
          lac i la
          jda lwt
          jsp lcc

vf2,      idx fa
          idx la
          sas rb1
          jmp vf1+2
          jsp rbk
          jmp vf1

ar,       dap arx      /assign reader
          law 51
          arq
          jmp bus
          jsp lct
          law 4642
          jda tys
          jsp soi
          jsp lct
arx,      jmp .

bus,      jsp lct      /busy
          lac (356224
          jda tys
          lac (223034
          jda tys
          jmp pn2

```


start

ID part 2 1/26/65.

lis,	law i 47 arq	/initialize assignments. /used by ff1
lis+2,	law i 51 arq	
zrt,	lac dpf dac dfp	/zero drum routine return
lse,	jsp lcc dzm mod dzm tas dzm iif	
lss,	clc dac chi	
lsp,	dzm wrd lac cun	
ssn,	dip sgn dzm dnm dzm syl	
n2,	dzm sym l dzm sym r clc dac let dac chc	
lsr,	dzm bbf lio skl dio wea init bax, lwt tyi	
ps1,	dio ch law dtb add ch dap .+1 lac .	
cas,	xx and (777 dac t2 sub (44 spa	/rar 9s or dio ch
cad,	jmp ln add (jmp uc dap lsx sub ar1 spq jmp i lsx lac sym l ior sym r lio lcf sad (45 dio iif law syl lio let spi i jsp ev1 jmp ev4	/used as add /last no-eval routine

law 77
sad ch
dzm tas
lac (flex U
jda tys
jmp lss

```

ev1,      dap evx  /symbol lookup
          law mst+2
          dap esk
evc,      lac est
          dap es4

evg,      dap ev2
ev2,      lac .
          spa
          jmp esn
          cla
es3,      sas sym l
          jmp esi
es4,      lac .
          sad sym r
          jmp ev3
esi,      idx es4
          idx es4
esk,      sub mst+2      /or (lac low
          sma
          jmp .+3
          add i esk
          jmp evg
          idx evx
ev3,      idx es4
evx,      jmp .
esn,      idx es4
          lac i ev2
          xor c4
          jmp es3

```

```

ev4,    dap sgn
        lac wrd
sgn,    xx /operator and syllable addr.
        dac wrd
        lio chi
        spi
        lac lwt
lsx,    jmp .

n,      rir 5s /number routine
        lac syl
        ral 3s
        spi i
cun,    ior t2
        dac syl
        lac dnm
        ral 2s
        add dnm
        ral 1s
        spi i
        add t2
        dac dnm
        jmp l1
ln,     add (44-12
        spa
        jmp n

        dzm let /letter routine
l1,     dzm chi
        idx t2
        idx chc
        sas (4
        jmp ln3
        lio sym r
        dio sym l
        dzm sym r
ln3,    sub (6
        szm
        jmp lsr
        lac sym r
        mul spd+1
        div spd+2
        jmp .
        add t2
        dac sym r
        jmp lsr

```

uc,n+44,	lio (rar 9s jmp .+2	/upper case
lc,	lio ps1 dio cas jmp lsr+1	/lower case
sqr,	lac dnm jmp n1+1	/' means take decimal number
quo,	lac sym l sza i lac sym r jda spv lac tys jmp n1	/" means take as flexo codes
q,	lac lwt jmp n1	/Q means last quantity
f,	lio chi spi i jmp ff law fg jmp n1-1	
a,	law ac jmp n1-1	/A means accumulator
ir,	law io jmp n1-1	/I means i-o
m,	law msk dac iif	/M means mask register
n1,	dzm chi dac syl jmp n2	
ta0, err, er1,	lac (743521 jda tys law 7234 jda tys law i 51 arq jmp lse	/? /lc, blk
ser,	law 7435 jda tys law 0772 jda tys lac est dap sr1	/symbol table overflow error

```
sr1,   lac  
       spa  
       idx sr1  
       idx sr1  
       lac i sr1  
       jda pi  
       lac (741034  
       jda tys  
       jmp lis+2
```

```
daq,   law 7777  
       and lwt  
       jmp .+2
```

```
com,   lac loc  
       dac df1
```

/D defines sym as address of Q

/comma defines sym as loc

```

def,      lac let          /define symbol
sk1,      sza
          jmp err
          law pn2

de,       dap dex
          lio df1
          jsp evl
          jmp df2
          law est-sl-sl
          sub est
          sma
          jmp ser          /symbol table full
          law i 1
          adm est
          dio i est
          sub one
          dap est
          lio sym r
          dio i est
          sub one
          swp
          lac sym l
          sza i
          jmp dex
          dio est
          xor c4
          dac i est
          jmp dex

df2,      dio i es4
dex,      jmp .
dot,      lio chc
          lac loc
          spi i
          lac dnm
          dac syl
          law 44
          dac t2
          jmp l1

del,      dzm iif          /end of no-eval routines, delete
          jmp pn2

```



```
val,      dac df1          /open paren, sets up value for define
          jmp lss

kil,      spi             /K
          jmp ki5
          lac let
          sza
          jmp err
lal,      lac low         /used below
          sad i est
          jmp err
          law ki1         /delete one symbol
          dap evx
          law lal
          jmp evl+2
```

eql,	jda eap jda opt	/print octal integer (=)
pn2,	jsp lct jmp lss	
arw,	jda eap	/print as instruction (→)
ar1,	jda pi jmp del	/used by cad+2
oct,	spi i jmp err law 10 jmp .+4	/octal-decimal switch setup /H
dec,	spi i jmp uns law 12 dap ops jmp lse	/U
smb,	spi i jmp sav law pi jmp .+2	/symbolic-constant switch setup /S
cns,	law opt dap pns jmp lse	/C
oad,	law pvl jmp t1s-1	/octal-relative switch setup /O
rad,	spi i jmp err law pev dap pa1	/R
t1s,	jmp lse	
pls,	lac cad jmp ssn	/plus, space
41n,	spi dio wrd lac csu jmp ssn	/minus
uni,	jmp ssn-1	/union, V
isc,	lac can jmp ssn	/intersection, ^

```

tab,      spi i           /tab
          jda dep
          dzm lcf

ta3,      dac lwt
          jsp lcc
          lio lcf
          sni
          jmp ta4
          sub (ac
          spa
          jmp ta0
          sub (nir       /internal symbol print
          sma
          jmp ta0
          add (nir
          sal 1s
          add (aa
          dap .+1 lac .
          jda tys
          idx .-2
          xct .-3
          jda tys
          jmp ta4+1

bs,       spi i           /backspace
          jda dep
          idx loc
          jmp ta3

fs,       spi i           /arrow up (forward space)
          jda dep
          law i 1
          adm loc
          jmp ta3

```

bac,	law opt jmp .+2	/open bracket (bar-constant)
bas,	law pi dap bax jmp bar	/closed bracket (bar-symbolic)
vb,	clc	/vertical bar
bar,	dac mod lac lcf dac tas lac iif dzm iif dac lcf sza dac tsf lac wrd spi 1 jmp ta5 lac tas dac lcf lac lwt dzm tsf jmp ta6	/slash . /tas used for temporary storage
uc8,	spi 1 jda dep dzm tsf jmp ta6	/> means make corr. and open register
cr,	spi 1 jda dep dac lwt law 72 jda tys jmp lse+1	/carriage return
eas,	law ws4+2 dap dpx law ea1 jmp ws	/effective address search
nws,	lac sk2 dac wea	/not word search
wds,	law ws1 jmp ws	/word search
pbx,	jda eap jda tys jmp pn2	/print as bcd (~)

```

rd,      spi i           /read binary tape (Y)
         jmp err
         jsp ar
         jmp .+2
nb,      jsp rbk
         law i 1
         add ch
         and (7777
         jda chk         /check last address of block
         lac ch
         sub t
         dap wc
         law 7777
         and t
         add dfp
         swp
         dia
         lio wc
         law buf
         dcc
         jmp dre
         jmp nb

```

```

bgn,    jmp bg0
xec,    jmp xe0
pra,    jmp pr0

```

```

bk,      spi i           /B
         jmp 1bk
         clc
         dac bbf
         dac let
         law bk1
         dac syl
         dac iif
         dzm chi
         jmp lsr+1
1bk,    dac dep
         law bk1
         add c4
         dac tas
         dac tsf
         jsp dep+1
         jmp lse

```

```

ovb,    clc           /overbar
         spi
         jmp n1
         sas bbf
         jmp err
         jsp cbp
         jmp lse

```

```

start

```

ID part 3 1/26/65.

```
t11,      spi 1          /title punch and punch format setup
          jmp err
          law 47
          arq
          jmp bus
          feed 30
          jsp lcc
          jmp tp1
```

```
jbk,      spi          /jump block
          jmp err
          law 47
          arq          /do we have the punch
          jmp err
          lac wrd
          ior cj
          dac lwt
          feed 40
          lio lwt
          jsp pbw
          feed 520
          jmp lis
```

```

pul,      jda chk          /lower limit setup
          dap fa
          jmp lss

pwd,      dac dep          /punch word
          lac tas
          sma
          jmp ta0
          lac tsf
          spi
          ior mod
          sza
          jmp ta0
          law 47
          arq          /do we have the punch?
          jmp bus
          spi
          jmp .+3
          jsp dep+1
          dac lwt
          cle
          dac pvf
          lac tas
          dap fa
          dap la
          jmp pb5

pun,      spi          /punch any length block
          jmp err
          law 47
          arq          /do we have the punch?
          jmp err
          jsp zro+1

pb5,      lac fa
          ior c77
          dac t
          sub la
          sma
          jmp pb6          /next hundred too high
          idx t

pb4,      jsp pbb          /pbb or pur
          lac t
          dap fa
          jmp pb5

pb6,      lac la
          dac t
          idx t
          xct pb4
          jmp pn2

```

```

vfy,      jsp zro+1          /for verify
          jmp ver

zro,      law zd           /for zero registers
          dap zvp         /used by block operations
          lac wrd
          spi
          jmp .+3
          jda chk
          jmp .+5
          cla
          dap fa
          law i 1
          add mb
          dac wrd
          dap la
          law 7777
          and fa
          jda chk
          dac chk
          dac lo
          sub wrd
          szm
          jmp err
          dzm dff
          dzm pvf
          law zd3
zvp,      jmp .

tbl,      spi i           /symbol table reader
          jmp tb2
          jsp ar
          lac t4
          sad (jmp 7750
          jmp tb5
          sad (jmp 6151
          jmp tb1-1
          sas (jmp 7751
          jmp err
          dzm sym 1

tbl,      jsp gwd         /reader-macro
          jda pz
          lac la
          sad rb1
          jmp tbn
          jsp gwd
          dac df1
          jsp de
          jmp tb1

```



```

tbn,    law est
        sub est
        sar 1s
        jda opt

tbm,    jsp rbk           /skips rest of tape
        jmp tbm

gwd,    dap gwx
        lac la
        sas rb1
        jmp gw1
        jsp rbk
        jmp gwd+1

gw1,    dap gwa
        idx la

gwa,    lac .
gwx,    jmp .

tb5,    jsp gwd           /read midas table
        dac sym 1
        jsp gwd
        and (177777
        dac sym r
        ior sym 1
        sza i
        jmp tbn
        jsp gwd
        dac df1
        lac sym 1
        lia
        ril 1s
        sma+spi-skp
        jmp tb5
        and (177777
        dac sym 1
        jsp de
        jmp tb5

```

tb2,	dzm sym 1	/table read from ts macro or possible
	sas one	
	jmp pot	/see if table read from possible
	lio mst	
	dia	
	law low	
	dac est	
	add dfp	
	cli+swp-opr	
	dcc	
	jmp dre	
	lac low-1	
	sub (lac 6150	
	sma	
	jmp err	
	add mst	
	spa	
	jmp err	
	law i 6151-low	
	add low-1	
	dac t3	
	law low-2	
tb3,	dap tb4	
tb4,	lac .	
	dac df1	
	law i 1	
	add tb4	
	dap tb4	
	lac i tb4	
	jda pz	
	jsp de	
	lac tb4	
	sad t3	
	jmp lse	
	sub (1	
	jmp tb3	

pot,	sas (2 jmp err lac (400006 dac tas jsp fet sal 1 cma dac opt idx tas jsp fet dac t1 lio t1 dia law 100 add opt sad (-100 jmp lse spa cla add (-100 dac t2 cma\lia adm opt lai adm t1 lai lor dfp lia law buf dap gsb dcc jmp dre	/not possible /twice no of symbols in possible /origin of table /finished /full block /word count(negative) /number of remaining words /initial drum address next /initial core address - reader buffer
gsb,	lac . and (177777 lia idx gsb dio sym r lac i gsb sni i jda df1 idx gsb idx t2 isp t2 jmp gsb jmp gfd	

```

pz,t6,      0                /permute zones. (temp storage here)
dap pzx
lac pz
and (202020
ral 1s
xor pz
xor c4
dzm sym r
lia
pz1,      cla
rcl 6s
add ps1+1
dap .+2
law 77
and .
dac t2
idx t2
lac sym r
ral 2s
add sym r
ral 3s
add t2
dac sym r
sni i
jmp pz1
pzx,      jmp .
ki1,      jmp .+2
          jmp err
          sub ev2
          dac t3
          idx t3
          dac t4
          idx t4
          idx es4
          and (7777
ki2,      sub t4
          dap ki3
          add t3
          dap ki4
ki3,      lio .
ki4,      dio .
          sas est
          jmp ki2
          add t3
          jmp .+2
ki5,      law low          /delete all symbols
          dac est
          jmp lse
fee,t2,    0                /feed subroutine and temp storage.
dap fex
cli
ppa
isp fee
jmp .-2
fex,      jmp .

```

```

lwt,      0          /Q, last word typed
          dap pnx
          lac lwt
pns,     jda pi      /pi or opt
pnx,     jmp .

eap,      0          /eql,arw,pbx common
          dap .+7
          lac eap
          dac lwt
          lac iif
          sza
          jmp ta0
          jsp let
          jmp .

ff,       lac sym l
          sza i
          lac sym r
          jda spv
          law ff1
          dap lsx
          lio let
          law syl
          spi i
          law tys
          jmp ev4      /returns with ac proper for arq
ff1,     lio io
          arq         /do the arq with phony ac but real io
          jmp ff2
          dio io      /F interacts only with io
          dac fft     /temp storage
ff3,     jsp lcc
          lac fft
          jda opt     /type out ac since it may contain info
          dac lwt
          jmp lse     /return to listen

/since lcc not transparent, save and restore ac etc

ff2,     dio io
          dac fft
          jmp ff3
fft,     0          /temp for ac

```

```

pi,t4,    xx                /print instruction
          dap px
          jsp pev
          lac pi
          sub ci
          spa
          jmp ppk
          dac pi
          law 72
          jda tys
          tyo
          law 71
          jda tys

ppk,      cli
          tyo
          szf 2
          jmp pvl
          law 72
          jda tys
          and (760000
          sad (sft
          jmp i66
          sad pr0+1          /law 0
          jmp plo
          rar 1s
          sza
csu,      sub (320000       /used as sub
          spa
          jmp plo
pvl,      lac pi
          sza i
          szf 1 i
pv3,      jda opt
px,       jmp .

i66,     law 1             /1s-9s
          add pi
          and pi
          sza
          jmp pvl
          law pa1+1
          dap pex
          lac ea
          jmp eak+2

pad,      0                /print address
          dap px
          law 7777
          and pad
          dac pi
          clf 1
pa1,     jsp pev          /pev or pvl
          lac (flexo +
          jda tys
          jmp pvl

```

ta4, jda pad
 lio mod
 law 7221
 spi
 law 7456 /for type-in mode
 jda tys

ta5, dzm loc
 dap loc
 lio lcf
 dio tsf

ta6, dap tas
 jsp lct
 lac cad
 dip tas
 lac mod
 sza

 jmp lss
 jsp fet
 dac lwt
bax, jda lwt /pi, opr or lwt
 jmp pn2

```

pev,      dap pex           /symbol lookup subr
          law i 7777
          and pi
          sad (opr         /detect operates
          jmp sev
          and (760000
          sad (skp         /detect skips
          jmp sev
          clf 2
eak,      lac est
          dap ea
          clf 1
eal,      lio i ea
          spi
          idx ea
          dap psw
          spi i
          cli
          dio t6
          idx ea
ea,        lac .
          szf 2           /test for "skip or operate" or other
          jmp sko
          xor pi
          spa
          jmp eix
          lac pi
          sub i ea
          spa
          jmp eix
          szf i 1
          jmp psw
          lac i ea
          sub i ei1
          szm
          jmp psw
eix,      idx ea
          lia
          sub mst+2
          swp
          spi i
          jmp .+3
          sas mst+1
          jmp eal
          szf 2
          jmp pex
          szf i 1
          jmp pvl
          lac pi
ei1,      sub
          lia
          sza           /detect neg nums
          jmp i77
          dio pi
eiy,      jsp spt
          lac pi
sk2,      sza i
          jmp px
          szf i 2

```



```

pex,      jmp .
          cma
          dac t2      /mask
          jmp eix

sev,      dac t1      /save instruction
          lac pi
          cma+stf 2-opr
          dac t2      /mask
          jmp eak

sko,      ior t1
          sas i ea
          jmp eix
          szf 1
          xor t1
          sza i
          jmp eix
          xor pi
          lia
          and t2
          sza
          jmp eix
          dio pi
          szf i 1
          jmp psw
          lac (flexo V
          jda tys

psw,      lac .      /best symbol thus far
          dac sym r
          lio t6
          dio sym 1
          lac ea
          dap ei1
          stf 1
          szf 2
          jmp eiy
          jmp eix

177,      law i 7777      /numbers 77xxxx
          and pi
          sas (770000
          jmp eiy-1
          law 7254
          jda tys
          lac pi
          cma
          jmp pv3

tys,      0      /type symbol
          dap tyx
          setup opt,3

```

```
tyl,   lac tys
        ral 6s
        dac tys
        and c77
        sza i
        jmp tyc

        sad (72
        jmp dns
        sad (74
        jmp ups
        swp
tyb,   tyo
tyc,   count opt, tyl
        lac lwt
        cli
tyx,   jmp .
```

```

dns,      lac ps1          /redundant case shift filter
          lio (72
dn1,      sad cas
          jmp tyc
          dac cas
          jmp tyb

ups,      lac (rar 9s
          lio (74
          jmp dn1

lcc,      dap lcx          /lower case, carriage return
          law 7277
          jmp lc1

lct,      dap lcx          /lower case, tab
          law 7236

lc1,      jda tys
lcx,      jmp .

so1,      rpb              /skip over input routine
soi,      rpb              /enter here
          spi i
          jmp so1
          dio t4

rbk,      dap rbx          /read a block into buffer
          init rb1, buf
          dap la
          dzm chi
          rpb
          dio t2
          dio t
          spi
          jmp lis+2        /start block read.
          rpb
          dio ch

          law i 1          /check for block format
          add ch
          sub t2
          and (777700
          sza
          jmp err

rb0,      rpb
rb1,      dio .
          lac i rb1
          adm chi
          idx rb1
          index t2, ch, rb0
          add chi
          add t
          rpb
          dio chi
          sad chi

rbx,      jmp .

```


/combined octal-decimal print subroutine

```
opt,      0
          dap opx
          dzm op1
opa,      lac opt
opb,      dac op2
          cli+swp-opr
          rcl 1s
          div opsops,      10
          sas op1
          jmp opb
          sni

          lio (20
          tyo
          lac op2
          dac op1
          sas opt
          jmp opa
opx,      jmp .
op1,      0

ci,       10000
c10,      10
```

/symbol print subroutine

```
spt,      dap spy
          lac sym 1
          and (177777
          jda spv
          jsp tys+1
          lac sym r
          jda spv
          jsp tys+1
spy,      jmp .
spv,      0
          dap spx
          init spj, spd
          dzm tys
spr,      lio spv
          cla
          rcl 1s
spj,      div
          .
          dio spv
          rar 1s
          add (spl-add
          dap .+1
          lio
          spa
          ril 6s
          lac tys
          rcl 6s
          dac tys
          idx spj
          sas (div spd+3
          jmp spr
spx,      jmp .
spd,      3100
          50
one,      1
spl,      flex 01
          flex 23
          flex 45
          flex 67
          flex 89
          flex ab
          flex cd
          flex ef
          flex gh
          flex ij
          flex kl
          flex mn
          flex op
          flex qr
          flex st
          flex uv
          flex wx
          flex yz
          737200
```

/dispatch table, LC, UC

dtb,

disp pls, pls
letter 1, quo
letter 2, sqo
letter 3, pbx
letter 4, daq
letter 5, uni
letter 6, isc
letter 7, pul

/space
/1,"
/2,'
/3,~
/4,∩
/5,∨
/6,∧

/7,<

letter 10, uc8
letter 11, fs
disp err, err
disp err, err
disp err, err
disp err, err
disp err, err
disp err, err

/8,>
/9,↑

letter 0, arw
disp bar, err
letter 34, smb
letter 35, tbl
letter 36, dec
letter 37, vfy
letter 40, wds
letter 41, xec

/0,→
//,?
/s
/t
/u
/v
/w
/x

letter 42, rd
letter 43, zro
disp err, err
disp com, eql
disp err, err
disp err, err
disp tab, tab
disp err, err

/y
/z
/,,=

/tab

disp pwd, err	/.,-
letter 23, jbk	/j
letter 24, kil	/k
letter 25, ttl	/l
letter 26, m	/m
letter 27, nws	/n
letter 30, oad	/o
letter 11, pra	/p
letter 32, q	/q
letter 33, rad	/r
disp err, err	
disp err, err	
disp min, pls	/-,+
disp def, bas	/),]
disp ovb, vb	/',]
disp val, bac	/([,]
disp err, err	
letter 12, a	/a
letter 13, bk	/b
letter 14, cns	/c
letter 15, pun	/d
letter 16, eas	/e
letter 17, f	/f
letter 20, bgn	/g
letter 21, oct	/h
letter 22, ir	/i
disp lc, lc	/lower case
disp dot, del	/.,x
disp uc, uc	/upper case.
disp bs, bs	/backspace
disp err, err	
disp cr, cr	/carriage return

start

ID part 4 4-22-65
/private variables.

wrd,	0		/quantity being assembled
sym,	0	0	/alpha symbol being assembled
chc,	0		/character count
chi,	0		/+0 when letter or number has been /typed since last typeout or c.r. input
let,	0		/+0 when letter in syllable, otherwise -0
ch,	0		/character
syl,	0		/syllable
t,	0		/temporary storage
la,	d10		/last address
fa,	d10		/first address
mod,	0		/mode, -0 for "type-in"
opd,dnm,	0		/decimal number
op2,	0		
/constants			
c4,	400000		
c77,	77		
c01,	010100		
cj,	600000		
c3,	3		
c2,	020000		
pvf,	.		/punch, verify flop /0 - usually, set by zro /-0 - center dot /not \pm 0 - continue in pv subroutine
bkf,	0		/breakpoint flop
esp,	dzm .		/used by zero routine
df1,	0		/value for defining symbol
loc,	0		/current location
tas,	0		/address part for fetch or deposit /current register. instruction part /+ for register closed, tells dep and /fet subroutine to ignore
tsf,	0		/current examination flop /0 - external register /non-0 - internal register
lcf,	0		/current location flop /0 - external /non-0 - internal
lif,	0		/initial internal flop /0 - usually /non-0 - set when type A,I,M,B
bbf,	0		/B flop /0 - usually /-0 - when B typed, not affected by uc,lc
cn3,	400000		/special proceed counter
dff,	0		/deposit-fetch flop, 0 dep, -0 fet
lo,	.		
wc,	.		

mst, 6151-low
lac tst+1-22
lac tst+1
630000

buf, 0 /reader buffer
t3, 0
t5, 0
buf+100/

pf, 0
pc, 0

ac, 0 /Internal registers.
io, 0
fg, 0
msk, -0
ll, 0
ul, 6277
bkl, -0
repeat nbp-1, -0

+.nbp/
bk0, -0 /switch for legal proceed
/if legal, has proceed address

awm, 30
t1, 0

c,

constants

start lis