

UNIVERSITY OF ILLINOIS
DIGITAL COMPUTER

LIBRARY ROUTINE M 20 - 234

By C. W. Gear

TITLE Eigenvalues of a Symetric Matrix by Givens' Method (SADOI Only)
 TYPE Automatic
 NUMBER OF WORDS 507 + Routines N12, R1, P16 and Y1
 MAXIMUM SIZE OF MATRIX 128 x 128
 DURATION Approximately $n^2/160 + n^3/20,000$ minutes
 ACCURACY Depends on matrix. So far as known, good to 11 decimal places for n less than 40 and good to 10 otherwise.
 SCALING The matrix must be scaled so that its norm is less than 1.

$$\sum_{ij} a_{ij}^2 < 1$$

CODED STOPS FF 036 in loc 322 on R.H.S. Sum check failed on read in.
 FF 037 in loc 1KK Norm of Matrix has exceeded 1.
 FF 038 in loc 074 Sum of squares check has failed. (see description)

FINAL STOP 24 014 on R.H.S. of loc 10N.

Routine is now ready to read in another matrix.

INPUT Normal input the program tape (it uses SADOI). It will stop on the word 34 011 FF 036 in loc 322 (16), on the L.H.S. if the tape has read in correctly, on the R.H.S. if there has been a reader error. Black switch the data tape in which should be written as follows:

Consider the matrix in its upper triangular form as

$$\begin{matrix} a_{11} & a_{12} & a_{13} & \dots & a_{1n} \\ & a_{22} & a_{23} & \dots & a_{2n} \\ & & & \dots & \end{matrix}$$

$$\begin{matrix} & & & & a_{n-1,n-1} & a_{n-1,n} \\ & & & & & a_{n,n} \end{matrix}$$

Each row should be punched, element by element as a sign followed by p decimal digits ($p \leq 12$) of the element, and

at the end of each row the character N should be punched. Thus the data tape will have the following sequence on it: n signed numbers, N, n-1 signed number, N, 3 signed numbers, N, 2 signed numbers, N, 1 signed number, N. The master routine will read in the first row of the matrix and plant n, the order of the matrix. It will then read in up to each N character, and store this as the next row of the triangular part of the matrix. If there are more than n-1 numbers in this row it will ignore all those after the n-1th, if there are less than n-1, say p, the p+1th will be the p+2th of the preceeding row, and so on to the n-1th interpreted as the nth of the preceeding row. It will stop reading in after it has sensed n N's.

MODIFICATION

This routine uses N12 to input the matrix. If it is desired to calculate the matrix elements within the Illiac prior to finding the eigenvalues the routine may be modified in the following way:

Replace N12 by an auxiliary subroutine to calculate the elements in the order in which they are read in. N12 is entered with the orders

$$\begin{array}{r}
 q \quad 50 \quad p \\
 \quad \quad 50 \quad q \\
 \hline
 \quad \quad 26 \quad 365F
 \end{array}$$

It occupies locations 365 through 403. These and the additional locations 670 through 869 may be used by the auxiliary subroutine. At no time must the auxiliary routine use other locations (except 0, 1, and 2.)

Successive rows of the upper triangular form should be calculated and stored in Williams memory locations p through 870+n-1. p is 870, 871, 872, ..., 870+n-1; on each successive entry. After the first entry to the auxiliary subroutine location 386 must contain 870 + n as a left hand address, where n is the order of the matrix.

The only elements affected in this transformation are those in the i th and j th rows and columns of \underline{A} . Since the matrix is symmetric, and this property is retained under an orthogonal rotation, we need only consider the columns of \underline{A} .

The elements of the rotated matrix \bar{A} are thus:

$$\bar{a}_{p,q} = a_{p,q} \quad p, q \neq i, j.$$

$$\bar{a}_{i,p} = \bar{a}_{p,i} = c_{ij} \cdot a_{pi} + s_{ij} \cdot a_{pj} \quad p \neq i, j.$$

$$\bar{a}_{j,p} = \bar{a}_{p,j} = c_{ij} \cdot a_{pj} + s_{ij} \cdot a_{pi} \quad p \neq i, j.$$

$$\bar{a}_{i,i} = c_{ij} (c_{ij} \cdot a_{ii} + s_{ij} \cdot a_{ij}) + s_{ij} (c_{ij} \cdot a_{ij} + s_{ij} \cdot a_{jj})$$

$$\bar{a}_{ij} = \bar{a}_{ji} = c_{ij} (c_{ij} \cdot a_{ij} + s_{ij} \cdot a_{ii}) + s_{ij} (c_{ij} \cdot a_{jj} - s_{ij} \cdot a_{ij})$$

Since the trace of a matrix is invariant under an orthogonal transformation,

$$\bar{a}_{jj} = a_{ii} + a_{jj} - \bar{a}_{ii}.$$

The values of c_{ij} and s_{ij} have not yet been assigned. In the Jacobi method, also in the library, they are chosen so as to make the (i,j) element zero. However, future rotations about other elements alter this, and iteration is necessary.

Suppose that rotations are made successively about the elements $(2,3):(2,4): \dots (2,n):(3,4):(3,5) \dots (3,n):(4,5) \dots (n-1,n)$; making at each step the elements $(1,3):(1,4): \dots (1,n):(2,4):(2,5): \dots (2,n):(3,5) \dots (n-2,n)$ zero. At any step, elements made zero in previous rotations are left zero.

Thus, all the elements of the matrix are reduced to zero except the diagonal and those adjacent to the diagonal.

This leaves the matrix

$$\begin{vmatrix} a_1 & b_1 & 0 & \dots & \dots & \dots & 0 \\ b_1 & a_2 & b_2 & & & & \\ 0 & b_2 & a_3 & b_3 & & & \\ & & & \dots & \dots & \dots & \\ & & & & b_{n-2} & a_{n-1} & b_{n-1} \\ 0 & \dots & \dots & \dots & 0 & b_{n-1} & a_n \end{vmatrix}$$

with eigenvalues equal to those of the original matrix.

Denote this matrix by B.

The eigenvalues of this are the roots of the polynomial

$$\det(B - uI) = f_n(u) = 0$$

From the form of the matrix the polynomial can be defined by the sequence

$$f_0(u) = 1$$

$$f_1(u) = a_1 - u$$

$$f_2(u) = (a_2 - u)f_1(u) - b_1^2 f_0(u)$$

.....

(1)

$$f_k(u) = (a_k - u)f_{k-1}(u) - b_{k-1}^2 f_{k-2}(u)$$

$$k = 2, 3, \dots, n.$$

If two successive f_i 's vanish then $f_n = 0$.

To overcome this, set $f_k = -b_{k-1}^2$ if $f_{k-1} = f_{k-2} = 0$ and $b_{k-1} \neq 0$ and $f_k = a_k - u$ if $f_{k-1} = b_{k-1} = 0$.

Then the number of roots of $f_n(x) = 0$ greater than u is

equal to the number of agreements in sign (zero considered as positive) in the sequence $f_0(u), f_1(u), f_2(u), \dots, f_n(u)$. For accuracy the equations are formed in semi-floating point.

b_i^2 is formed and scaled so that:

$$0 \leq b_i^2 \cdot 2^{-\epsilon_i} q_i < 2^{-40}$$

$$\text{or } q_i = 0 \quad \text{and } \epsilon_i = 79 \quad \text{if } b_i \neq 0$$

$$\frac{1}{2} \leq q_i < 1$$

and $a_i - u$ is scaled so that:

$$a_i - u = p_i 2^{-\delta_i}, \quad 0 \leq \delta_i \leq 39; \quad \frac{1}{2} \leq p_i < 1$$

$$\text{and } f_i(u) = 2^{-\pi_i} n_i$$

If π_0 is set as 0, and $v_i = \pi_i - \pi_{i-1}$ the machine equations are in the form:

$$m_i = 2^{\sum_i} (p_i m_{i-1} - 2^{\Delta_i} q_{i-1} m_{i-2}) \quad \text{if } \Delta_i < 0$$

$$\sum_i = v_i - \delta_i$$

$$= 2^{\sum_i} (2^{-\Delta_i} p_i m_{i-1} - q_i m_{i-2}) \quad \text{if } \Delta_i \geq 0$$

$$\sum_i = v_i + v_{i-1} - \epsilon_{i-1}$$

$$\text{where } \Delta_i = v_{i+1} + \delta_i - \epsilon_{i-1}$$

\sum_i is found by scaling the quantity in parenthesis to lie between $1/2$ and 1 , and v_i is found from the appropriate equation.

THE CODE

The program is read into the memory and stored on the drum for future use as temporary storage overwrites much of it. The triangular part of the matrix is read in row by row and packed on the drum according to the following arrangement: Consider a track of length 65 words,

location	+ 0	1	2	3 n-64	64
2560	a_{11}	0	$a_{66,66}$	$a_{66,67}$	$a_{66,n}$	0
2625	a_{12}	a_{22}	0	$a_{67,67}$	$a_{67,n}$	0
2690	a_{13}	a_{23}	a_{33}	0		
.....
2560 + 64x65	$a_{1,65}$	$a_{2,65}$	$a_{65,65}$
2560 + 65x65	$a_{1,66}$	$a_{2,66}$	$a_{65,66}$
.....
2560 + (n-1)65	$a_{1,n}$	$a_{2,n}$	$a_{65,n}$

If two locations A and B are such that $\text{loc}(B) - \text{loc}(A) = 1 \pmod{64}$ then B can be utilised in the minimum access time after A has been referred to.

If p is greater than 65, then $a_{65,p}$ is in location $2560 + 64 \times 66 + 65(p - 65)$, and $a_{66,p}$ is in location $p - 64 + 2560$.

$$\begin{aligned}
 \text{Therefore } \text{loc}(a_{66,p}) - \text{loc}(a_{65,p}) &\pmod{64} \\
 &= -64 \times 66 - 65(p-65) + p - 64 \pmod{64} \\
 &= 65^2 - 64p \pmod{64} \\
 &= (64 + 1)^2 \pmod{64} \\
 &= 1 \pmod{64}
 \end{aligned}$$

Therefore the rows and columns of the matrix can be referred to in minimum access time.

The sum of the squares of the elements of the matrix is calculated as it is read in, and stored as a check against the sum after the reduction to Jacobi form.

After this reduction, the program will hang up if the difference exceeds 100×2^{-39} .

The eigenvalues are then found from the reduced form, by storing the square root R of the sum of squares of the original matrix as an upper and lower bound on the absolute size of the eigenvalues and adjusting these until they converge with information gained about the position of the roots by the method outlined.

Taking $P(u)$ to be the number of eigenvalues greater than u , the method can best be seen from a flow chart.

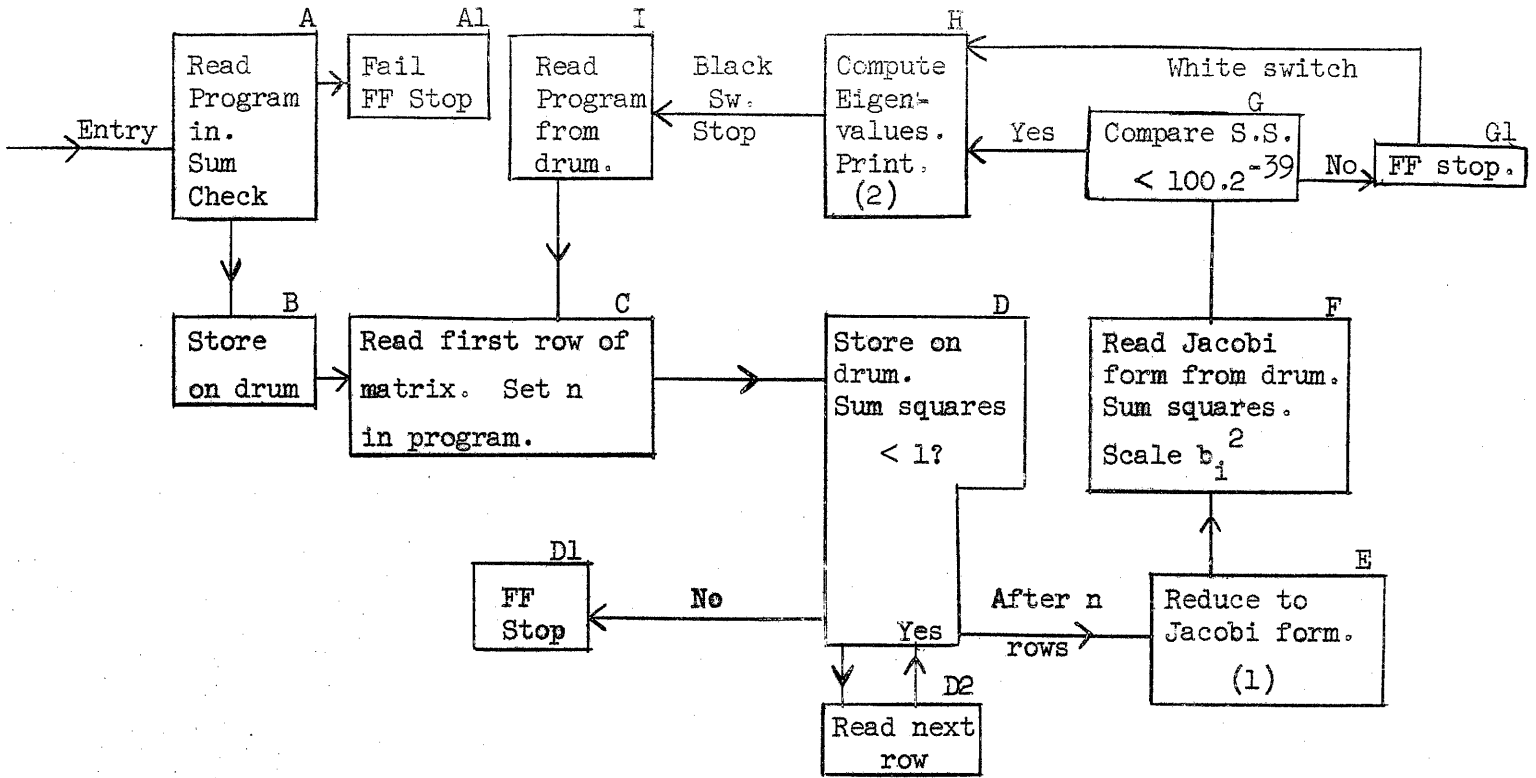
STORAGE ALLOCATION

The program occupies locations 17 through 668 and also on the drum. During input of the matrix locations 870 - 997 are used to store the rows. During reduction to Jacobi form locations 742 - 869 and 870 - 997 are used as temporary storage for the $i+1$ and $j+1$ th rows of the matrix which are altered in a rotation about the $i+1, j+1$ element. During the calculation of the eigenvalues the a_i 's are stored in 742-869, the scaled b_i^2 in locations 870 - 997 with their scaling factors in 614 - 741. Locations 358 - 485 hold the lower bounds of the eigenvalues, while locations 486 - 613 hold the upper values.

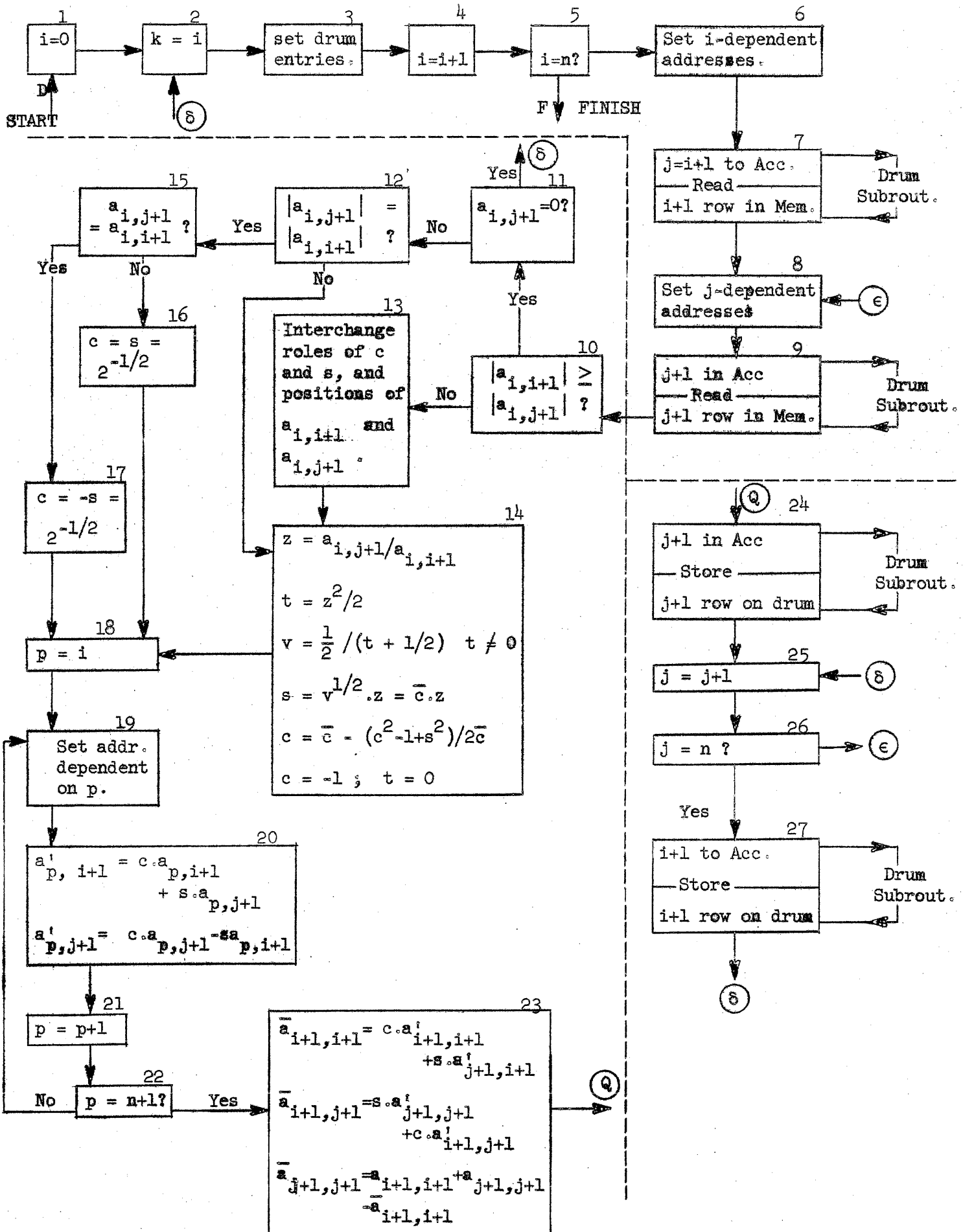
DATE June 27, 1957 RT:1/23/59

PROGRAMMED BY *B. W. Gear*

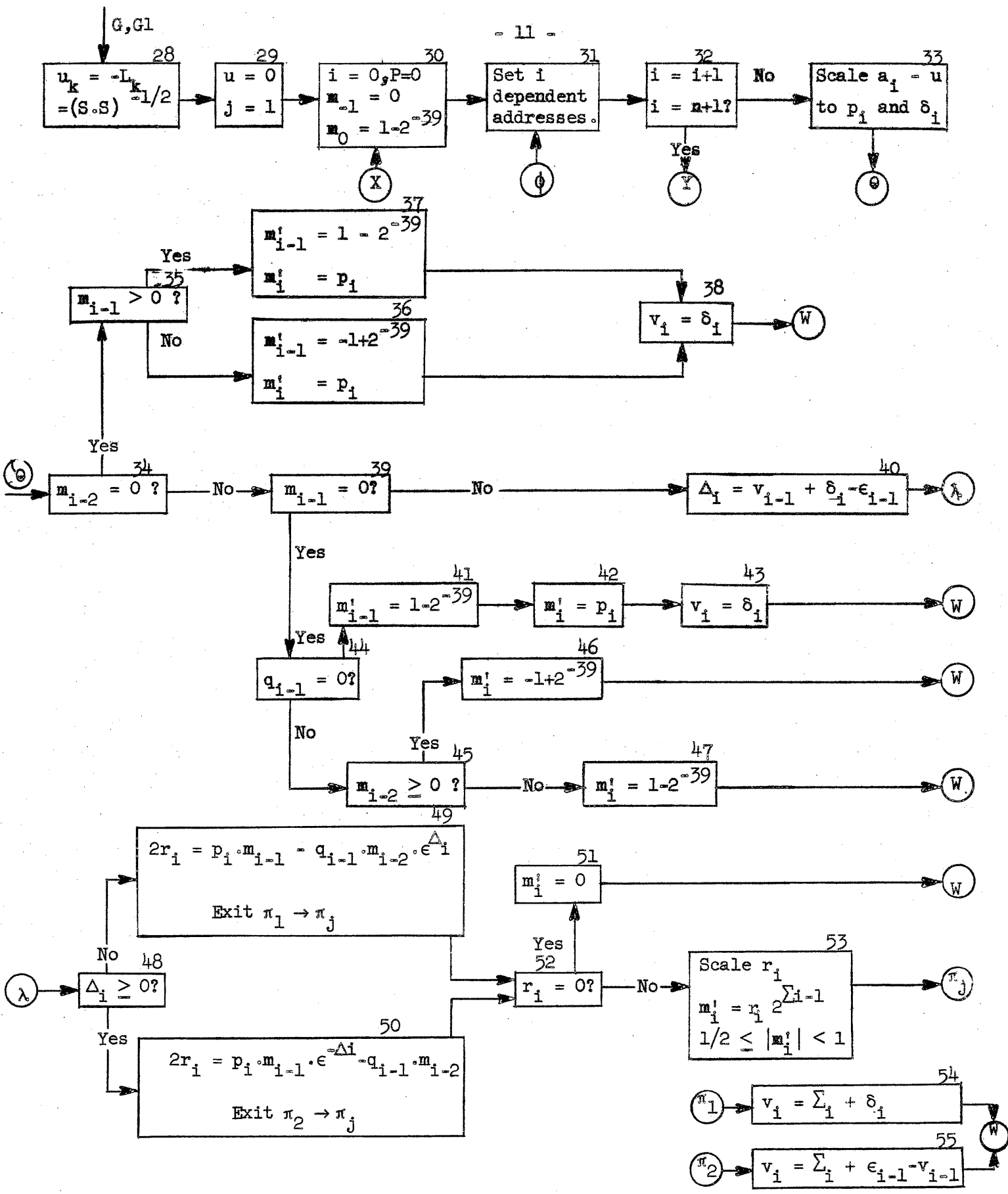
APPROVED BY *D. E. Muller*



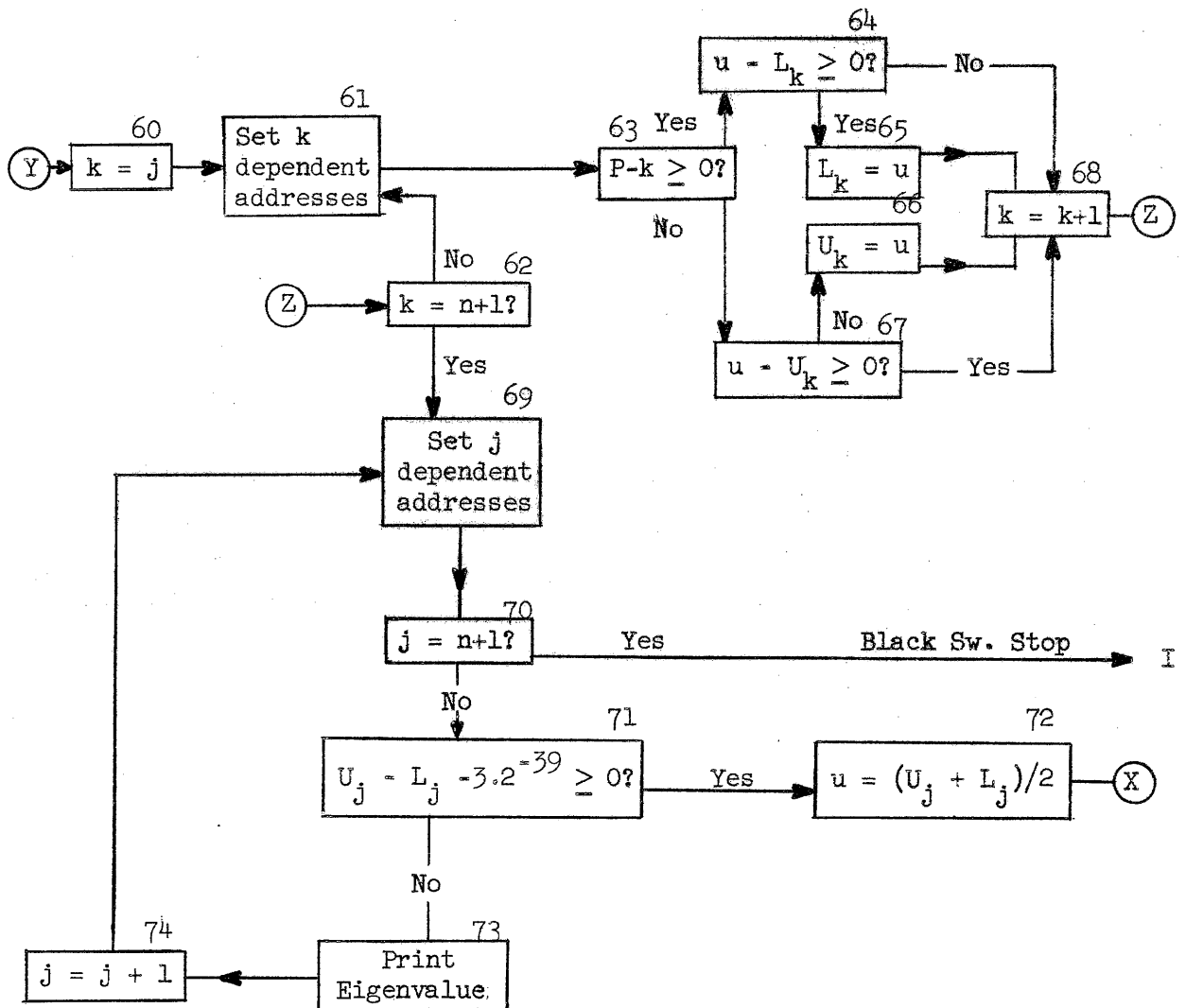
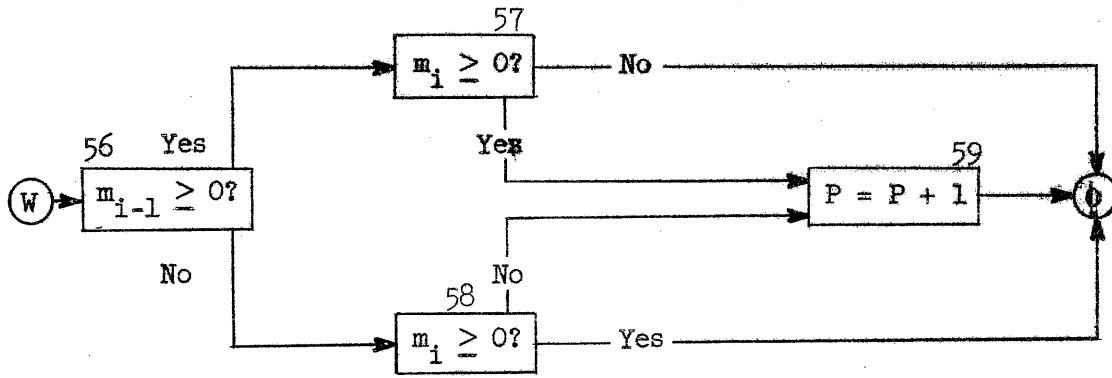
MAIN FLOW CHART



(1) Reduction to Jacobi Form



(2) Eigenvalues. Part I.



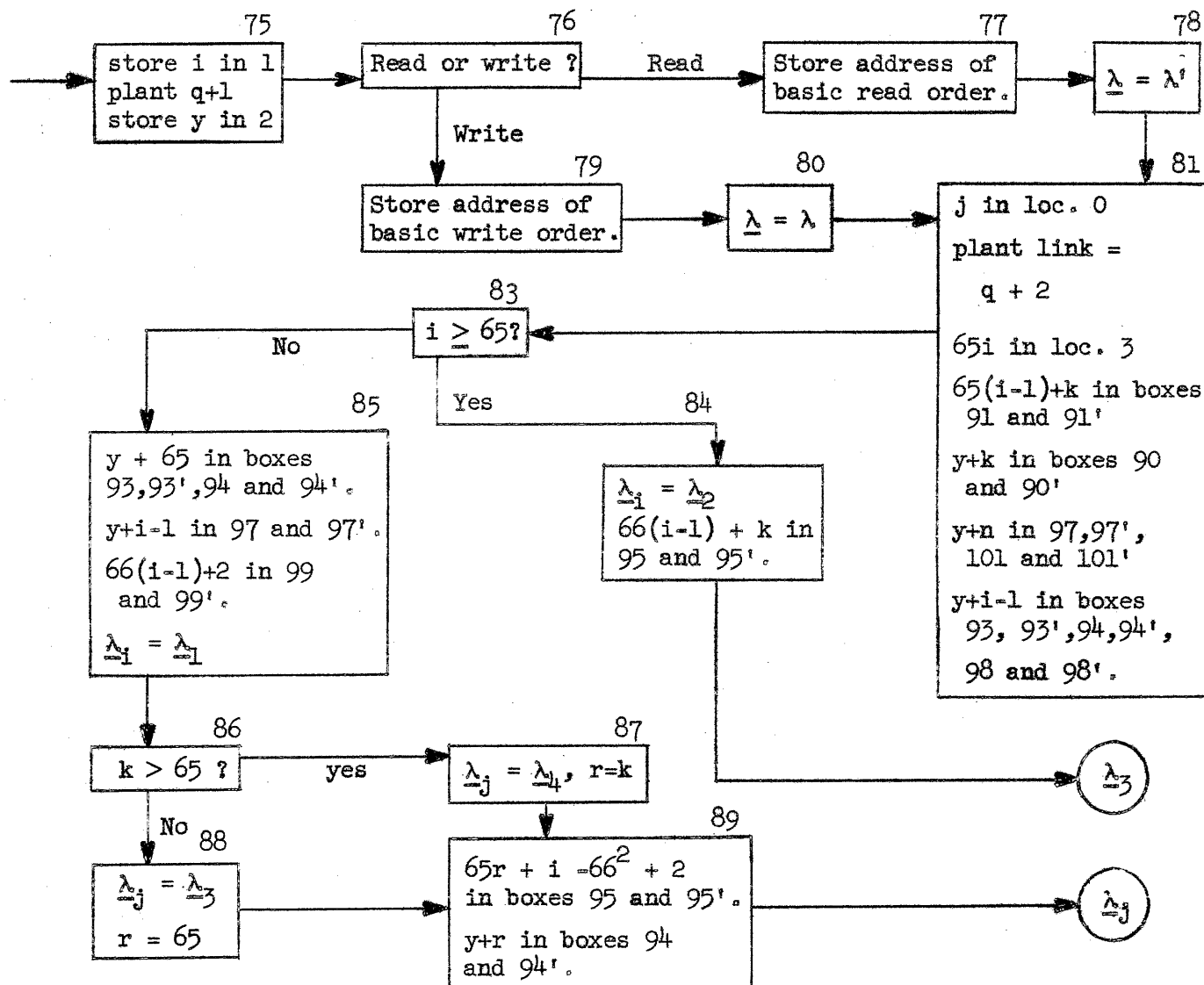
(2) Eigenvalues. Part II.

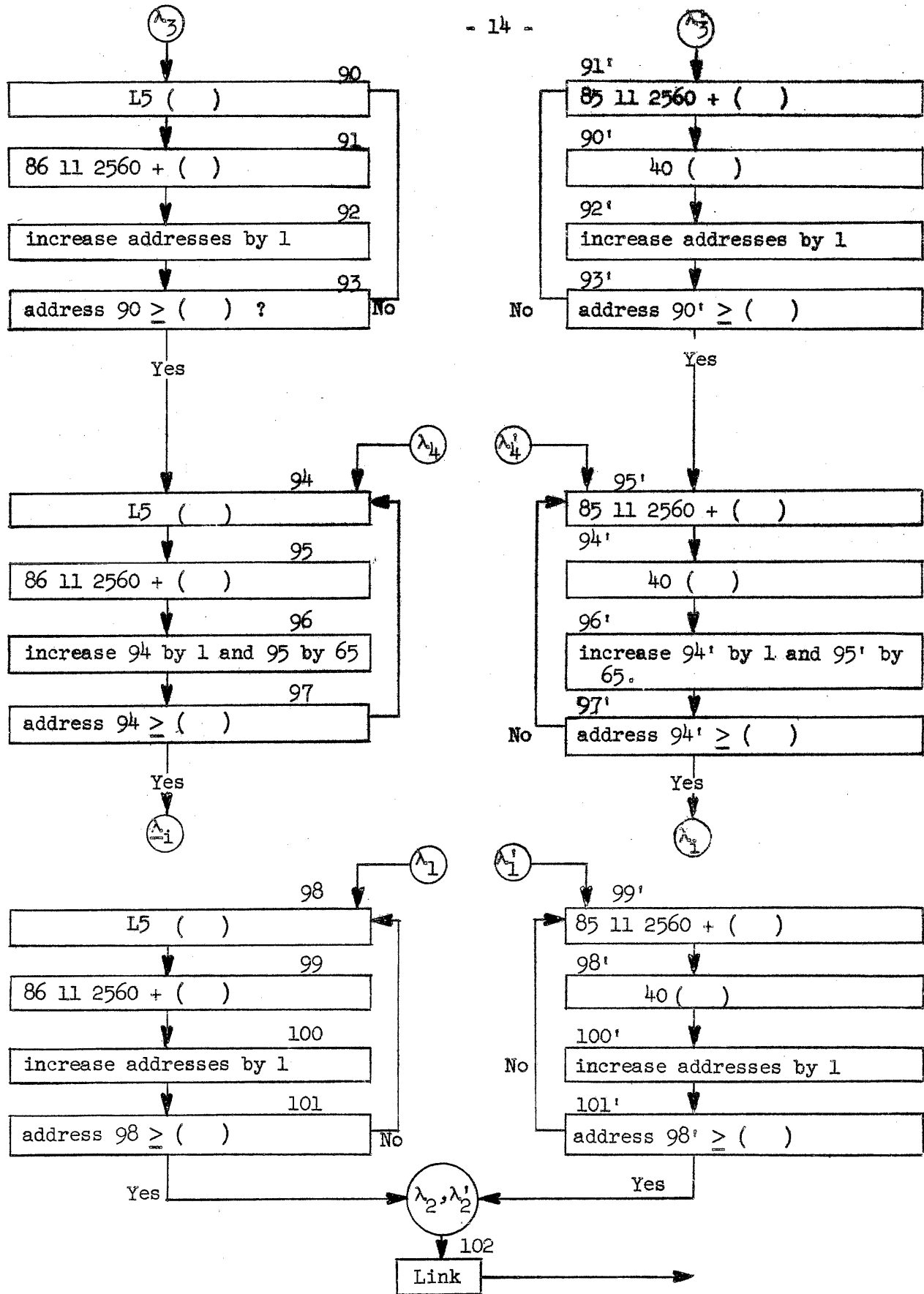
The Drum Subroutine

This is an independent subroutine which stores the triangular part of a symmetric matrix on the drum as outlined above. Entry is made to read (write) the i th row except for the first j elements into (from) Williams memory locations $y+j$ through $y+n$ with the following orders:

q	$5(J)0$	y
	50	q
$q+1$	26	--
	00	k

with i in the accumulator.





The Drum Subroutine. Part II.

LOCATION	ORDER	BOX NO.	NOTES	PAGE 1
17	KO 017K JO 100F 50 17F	B	Store program on drum.	
18	26 23F 02 2688F			
19	00 840F 26 404F		Exit to C.	
20	50 100F 50 20F	I	Read program from drum.	
21	26 23F 02 2688F			
22	00 840F 26 404F (Y1) 00K		Exit to C.	
63	F5 339F 42 84F		Prepare to read Jacobi form from drum.	
64	42 85F 42 72F			
65	F4 341F 42 91F			
66	42 105F 42 78F			
67	42 92F F5 84F			
68	L0 341F 42 107F			
69	L5 336F 40 71F			
70	L4 334F 40 77F			
71	00 0F 00 0F		Read from drum locations 2560, 2625, 2626, 2690, 2691, ...	
72	22 72F 40 0F			

LOCATION	ORDER	BOX NO.	NOTES	PAGE 2
73	L5 71F F4 334F		After 130 elements have been read reset loop	
74	40 71F F4 72F		and read 2562, 2563, 2628, 2629, ... to end of	
75	42 72F L0 347F		matrix.	
76	36 83F 26 77F		Test for end.	
77	00 OF 00 OF			
78	22 78F 40 OF			
79	L5 77F F4 334F			
80	40 77F F5 78F			
81	42 78F L0 662F			
82	36 71F 26 663F		Test for 130 elements.	
83	41 10F 41 11F		Prepare to sum squares and to scale b_i^2 's.	
84	41 12F 50 OF			
85	L5 11F 74 OF		Sum squares of a_i 's.	
86	L4 10F 40 10F			
87	S5 OF 40 11F			
88	F5 84F 42 84F			
89	42 85F L0 346F			

LOCATION	ORDER	BOX NO.	NOTES
90	32 84F 41 13F		
91	22 91F 50 OF		Square b_i 's and scale.
92	22 92F 75 OF		Also sum squares.
93	40 OF S5 OF		
94	40 1F L4 13F		
95	10 39F 40 14F		
96	L7 14F L4 OF		
97	L4 12F 40 12F		
98	S5 OF 40 13F		
99	L3 OF 36 128F		
100	41 2F 50 1F		
101	L5 OF 00 1F		
102	40 OF F5 2F		
103	42 2F L5 OF		
104	32 101F 10 1F		
105	L4 338F 40 OF		
106	L5 2F F0 335F		

LOCATION	ORDER	BOX NO.	NOTES	PAGE 4
107	22 107F 40 OF			
108	F5 107F 42 107F			
109	F5 91F 42 91F			
110	42 92F 42 105F			
111	L0 345F 32 91F			
112	L5 10F L4 12F		Sum a_i^2 + twice sum b_i^2 - previous sum of squares.	
113	L4 12F L0 351F			
114	40 10F L7 354F			
115	L2 10F 32 116F			
116	FF 56F L5 127F		Test to see if it exceeds 100.2^{-39}	
117	40 614F 41 870F		Set $b_0^2 = 0$, scaling factor = 79.	
118	F5 269F 42 120F			
119	F5 270F 42 121F			
120	L5 353F 40 OF	28	Set upper bound for eigenvalues = square root of sum of squares, and lower bound = - square root.	
121	L1 353F 40 OF			
122	F5 120F 42 120F			
123	F5 121F 42 121F			

LOCATION	ORDER	BOX NO.	NOTES	PAGE 5
124	L0 274F 36 120F			
125	41 273F F5 273F	29	$u = 0$	
126	40 272F 26 131F		$j = 1$	
127	00 0F 00 79F		Const.	
128	L3 1F 32 129F		If $b_i^2 = 0$ set scaling factor = 79.	
129	26 100F L5 127F			
130	40 2F 23 105F			
131	41 221F 41 5F	30	$i = 0$ $m_{-1} = 0$	
132	41 222F 41 3F			
133	L5 352F 40 4F			
134	F5 222F 42 222F	31	$M_0 = 1 - 2^{-39}, i = 1$	
135	L4 339F 42 142F			
136	L0 341F 42 161F			
137	42 198F L4 341F		Set i dependent addresses	
138	L4 341F 42 158F			
139	42 206F L0 349F			
140	36 231F L5 273F	32	$i = n+1 ?$	

LOCATION	ORDER	BOX NO.	NOTES	PAGE 6
141	10 1F 40 6F	33	Scale a_i-u	
142	43 150F L5 0F			
143	10 1F L0 6F			
144	40 1F L3 1F			
145	36 177F L5 1F			
146	40 6F F5 150F			
147	42 150F L7 6F			
148	00 1F 36 146F			
149	L5 1F 50 335F			
150	10 1F 00 0F			
151	40 1F L5 150F			
152	L0 275F 40 2F			
153	L3 3F 36 216F	34	$m_{i-2} = 0?$	
154	L3 4F 32 206F	39	$m_{i-1} = 0?$	
155	51 1F 74 4F		Form $q_{i-1} \cdot m_{i-2}$. Store double precision in 6 and 11	
156	10 1F 40 6F			
157	S5 0F 40 11F			

LOCATION	ORDER	BOX NO.	NOTES
158	51 3F 74 0F		Form $p_i \cdot m_{i-1}$. Store double precision in 7 and 12
159	10 1F 40 7F		
160	S5 0F 40 12F		
161	L5 5F L0 0F	40	$\Delta_i = v_{i-1} + \delta_i + \epsilon_{i-1}$
162	L4 2F 40 8F		
163	43 193F L7 8F		
164	42 226F 42 227F		Store as scaling factor
165	L0 230F 36 225F		$\Delta_i \geq 40 ?$
166	L5 227F 40 171F		
167	40 182F L5 8F		
168	36 179F L5 229F	48	$\Delta_i \geq 0 ?$
169	46 198F L1 12F		
170	10 39F L0 7F	49	Calculate r_i if $\Delta_i < 0$
171	00 0F 00 0F		
172	L4 6F 40 6F		
173	S5 0F L4 11F		
174	10 39F 40 0F		

LOCATION	ORDER	BOX NO.	NOTES	PAGE 8
175	L1 OF L4 6F			
176	40 6F 26 186F			
177	L5 276F 40 2F			
178	41 1F 26 153F			
179	L5 228F 46 198F			
180	F5 182F 42 182F			
181	50 11F L5 6F	50	Calculate r_i if $\Delta_i \geq 0$	
182	00 OF 00 OF			
183	00 1F L0 7F			
184	40 6F 85 OF			
185	L0 12F 26 174F			
186	L3 6F 36 223F	52	$r_i = 0 ?$	
187	S5 OF 40 OF			
188	L5 6F 40 1F	53	Scale r_i	
189	F5 193F 42 193F			
190	L7 1F 00 1F			
191	40 1F 36 189F			

LOCATION	ORDER	BOX NO.	NOTES	PAGE 9
192	L5 6F 50 OF			
193	10 1F 00 OF			
194	40 3F 50 3F			
195	L5 4F 40 3F			
196	S5 OF 40 4F			
197	L5 193F L0 275F			
198	22 OF L4 OF			
199	L0 5F 40 5F	55	$v_i = \sum_i + \epsilon_{i-1} = v_{i-1}$	
200	22 201F L4 2F			
201	40 5F L5 3F	54 56	$v_i = \sum_i + \delta_i$ $m_{i-1} \geq 0 ?$	
202	36 205F L5 4F	58	$m_i \geq 0 ?$	
203	36 134F F5 221F	59	$P = P + 1$	
204	42 221F 26 134F			
205	L5 4F 32 203F	57	$m_i \geq 0 ?$	
206	26 134F L3 OF	44	$q_{i-1} = 0 ?$	
207	32 212F L5 3F			
208	32 210F 41 3F	45	$m_{i-2} \geq 0 ?$	

LOCATION	ORDER	BOX NO.	NOTES
209	L5 352F 40 4F	47	$m_i^i = 1-2^{-39}$
210	22 201F 41 3F		
211	L1 352F 40 4F	46	$m_i^i = -1+2^{-39}$
212	22 201F L5 1F	42	$m_i^i = p_i$
213	40 4F L5 352F	41	$m_{i-1}^i = 1-2^{-39}$
214	40 3F L5 2F	43	$v_i = \delta_i$
215	40 5F 22 201F		
216	L1 4F 32 217F	35	$m_{i-1}^i > 0 ?$
217	22 212F L1 352F		
218	40 3F L1 1F		
219	40 4F L5 2F	36 38	$m_{i-1}^i = -1+2^{-39}; m_i^i = -p_i$ $v_i = \delta_i$
220	40 5F 22 201F		
221	00 0F 00 0F		P
222	00 0F 00 0F		Counter for sturm sequence. (i)
223	L5 4F 40 3F	51	$m_i^i = 0$
224	41 4F 22 201F		
225	L5 226F 22 166F		If scaling factor of Δ_i exceeds 40 avoid subtraction in boxes 49 and 50

LOCATION	ORDER	BOX NO.	NOTES	PAGE 11
226	41 335F 50 OF			
227	40 OF 10 OF		Constants	
228	00 198F 00 OF			
229	00 200F 00 OF			
230	00 OF 00 45F			
231	15 272F 40 271F	60	$K = j$	
232	15 271F 14 269F	61	Set K dependent addresses.	
233	42 237F 42 262F			
234	15 271F 14 270F			
235	42 263F 42 266F			
236	15 221F 10 271F	63	$P - k \geq 0 ?$	
237	32 263F 15 OF	67	$u - u_k \geq 0 ?$	
238	F0 273F 36 262F			
239	10 1F F5 271F	68	$k = k + 1$	
240	40 271F 10 348F	62	$k = n + 1 ?$	
241	36 232F 15 272F			
242	14 270F 42 246F	69	Set j dependent addresses	

LOCATION	ORDER	BOX NO.	NOTES
243	42 257F L5 272F		
244	L4 269F 42 245F		
245	42 259F L5 OF		
246	L0 355F L0 OF	71	$U_j - L_j - 3.2^{-39} \geq 0 ?$
247	32 257F 92 131F	73	Print eigenvalues and number of root (j).
248	92 515F 92 967F		
249	L5 272F 22 250F		
250	J2 3F 50 250F		
251	26 277F 92 67F		
252	92 963F L5 273F		
253	50 12F 50 253F		
254	26 277F F5 272F		
255	40 272F L0 348F	74	$j = j + 1$
256	32 241F 22 268F	70	$j = n + 1 ?$
257	00 OF L5 OF	72	$u = (U_j + L_j)/2$
258	10 1F 40 OF		
259	22 259F L5 OF		

LOCATION	ORDER	BOX NO.	NOTES
260	10 1F 14 OF		
261	40 273F 26 131F		
262	15 273F 40 OF	66	$U_k = u$
263	22 239F 11 OF		
264	14 273F 36 266F	64	$u - L_k \geq 0 ?$
265	22 239F 00 OF		
266	15 273F 40 OF	65	$L_k = u$
267	22 239F 00 OF		
268	22 268F 24 20F		Useful constants
269	00 OF 00 485F		
270	00 OF 00 357F		
271	00 OF 00 OF		
272	00 OF 00 OF		
273	00 OF 00 OF		
274	71 353F 40 486F		
275	10 1F 00 2F		
276	00 OF 00 79F		

LOCATION	ORDER	BOX NO.	NOTES	PAGE 14
	(P16) OOK			
333	00 OF		More useful constants	
	00 OF			
334	00 OF			
	00 65F			
335	00 OF			
	00 OF			
336	85 11F			
	00 2560F			
337	86 11F			
	00 2560F			
338	80 OF			
	00 OF			
339	00 OF			
	00 741F			
340	50 OF			
	00 OF			
341	00 128F			
	00 128F			
342	00 OF			
	00 OF			
343	80 OF			
	00 OF			
344	40 1F			
	50 OF			
345	K2 91F			
	50 OF			
346	M1 12F			
	50 OF			
347	22 72F			
	40 OF			
348	80 OF			
	00 OF			

LOCATION	ORDER	BOX NO.	NOTES
349	00 128F 00 OF		
350	50 OF 26 424F		
351	00 OF 00 OF		
352	7L 4095F LL 4095F		1 - 2 -39
353	00 OF 00 OF		
354	00 OF 00 100F		
355	00 OF 00 5F		
356	(R1) 00K		
365	(N12) 00K		
404	41 9F 41 10F		Program starts here
405	41 4F 22 406F		
406	50 870F 50 406F		
407	26 365F L5 386F		Read in the first row into locations 870 - 870+n-1.
408	46 350F L0 341F		870 + n in location (N12) + 21.
409	46 340F 10 20F		Set the test constants dependent on n.
410	42 347F 42 346F		
411	L4 341F 42 349F		

LOCATION	ORDER	BOX NO.	NOTES
412	42 345F L0 341F		
413	L0 339F 42 348F		
414	F0 335F 42 343F		
415	42 333F F0 335F		
416	42 342F L5 446F		and arrange for program to skip this in future
417	40 407F L5 4F		
418	42 420F F5 4F		
419	J0 870F 50 419F		Store using drum subroutine
420	26 555F 00 0F		
421	L5 447F 40 424F		
422	L5 406F 46 423F		
423	50 0F 46 424F		Sum squares, taking twice the sum of the off diagonal elements.
424	75 0F 00 0F		
425	L4 9F 40 9F		
426	36 427F FF 55F		Stop if norm exceeds 1.
427	85 0F L4 10F		
428	10 39F 40 0F		

LOCATION	ORDER	BOX NO.	NOTES
429	L7 OF L4 9F		
430	40 9F S5 OF		
431	40 10F L5 448F		
432	40 424F L5 423F		
433	L4 445F 46 423F		
434	L0 350F 36 436F		
435	L5 423F 26 423F		
436	40 10F L5 406F		Increase addresses and test.
437	L4 445F 46 406F		
438	F5 4F 42 4F		
439	L0 343F 32 406F		
440	L5 443F 40 407F		Reset routine
441	L5 444F 40 406F		
442	41 OF 26 449F		
443	26 365F L5 386F		Constants
444	50 870F 50 406F		
445	00 1F 00 0F		

LOCATION	ORDER	BOX NO.	NOTES	PAGE 18
446	26 365F			
	22 417F			
447	75 0F			
	26 425F			
448	75 0F			
	00 1F			
449	15 9F		Compute square root of sum of squares.	
	50 449F			
450	26 356F		<u>Jacobi reduction</u>	
	40 353F			
451	15 9F			
	40 351F			
452	43 554F	1	i = 0	
	15 554F		k = i	
453	42 468F	3	Set drum subroutine entries	
	42 475F			
454	42 543F			
	42 548F			
455	F5 554F	4	i = i+1	
	42 554F			
456	L0 342F	5	i = n ?	
	36 63F			
457	L5 339F	6	Set i dependent addresses	
	L4 554F			
458	42 476F			
	L4 341F			
459	42 477F			
	42 344F			
460	00 20F			
	46 541F			
461	F5 476F			
	00 20F			
462	46 532F			
	46 535F			

LOCATION	ORDER	BOX NO.	NOTES	PAGE 19
463	46 540F			
	46 516F			
464	14 341F			
	46 539F			
465	46 536F			
	46 515F			
466	F5 554F	7	$j = i+1$	
	42 552F			
467	50 742F			
	50 467F			
468	26 555F		Enter drum subroutine to read $i+1$ th row.	
	00 0F			
469	F5 339F	8	Set j dependent address	
	14 552F			
470	42 514F			
	42 533F			
471	42 538F			
	14 341F			
472	42 515F			
	42 537F			
473	42 540F			
	F5 552F	9	Enter drum subroutine to read $(j+1)$ th row	
474	50 870F			
	50 474F			
475	26 555F			
	00 0F			
476	41 0F			
	15 0F			
477	40 3F			
	15 0F			
478	40 4F			
	13 4F	11	$a_{1,j+1} = 0?$	
479	36 544F			
	L7 3F	10	$ a_{1,j+1} \geq a_{1,j+1} ?$	

LOCATION	ORDER	BOX NO.	NOTES
480	L2 4F 36 506F		
481	50 3F L5 4F	13	Interchange roles of c and s and positions of $a_{i,i+1}$ and $a_{i,j+1}$
482	40 3F S5 0F		
483	40 4F L5 551F		
484	42 497F 46 504F		
485	50 335F L5 4F	14	Calculate c and s
486	66 3F S5 0F		
487	40 5F 7J 5F		
488	10 1F 40 1F		
489	L3 1F 36 505F		
490	LJ 1F 40 1F		
491	50 335F 6L 1F		
492	K5 0F 36 494F		
493	26 505F 00 0F		
494	S5 0F 50 494F		
495	26 356F 40 1F		
496	50 1F 7J 5F		

LOCATION	ORDER	BOX NO.	NOTES
497	40 OF 40 OF		
498	50 OF 79 OF		
499	40 OF 50 1F		
500	79 1F L0 338F		
501	L4 OF 50 335F		
502	66 1F 85 OF		
503	10 1F L4 1F		
504	40 OF 22 514F		
505	L5 338F 22 495F		
506	40 1F L3 1F	12	$ a_{i,j+1} = a_{i,i+1} ?$
507	36 509F 26 508F		
508	L5 550F 26 484F		
509	L5 3F L4 4F	15	$a_{i,j+1} = a_{i,i+1} ?$
510	40 OF L5 553F		
511	40 5F 40 6F	16	$c = s = 2^{-\frac{1}{2}}$
512	L3 OF 32 513F		
513	22 514F L1 553F	17	$c = -s = 2^{-\frac{1}{2}}$

LOCATION	ORDER	BOX NO.	NOTES
514	40 6F L5 0F		
515	40 0F L5 0F		
516	L4 0F 40 0F		
517	L5 344F 42 526F	18	p = 1
518	42 530F 42 524F	19	Set p dependent addresses
519	L0 341F 42 529F		
520	00 20F 46 528F		
521	46 523F L0 340F		
522	36 532F 26 523F	22	p = n+1 ?
523	50 0F 7J 6F	20	Rotate matrix, multiplication loop
524	40 1F 50 0F		
525	7J 5F L0 1F		
526	40 1F 50 0F		
527	7J 6F 40 2F		
528	50 0F 7J 5F		
529	L4 2F 40 0F		
530	L5 1F 40 0F		

LOCATION	ORDER	BOX NO.	NOTES
531	F5 526F 22 517F	21	P → P+1
532	50 OF 7J 5F	23	
533	40 LF 50 OF		
534	7J 6F 14 LF		
535	40 OF 50 5F		
536	7J OF 40 LF		
537	50 6F 7J OF		
538	14 LF 40 OF		
539	40 OF 15 OF		
540	10 OF 40 OF		
541	41 OF F5 552F	24	Store j+1 th row on drum
542	J0 870F 50 542F		
543	26 555F 00 OF		
544	F5 552F 42 552F	25	j = j+1
545	10 343F 36 469F	26	j = n ?
546	F5 554F 22 547F		
547	J0 742F 50 547F	27	Store i+1 th on drum

LOCATION	ORDER	BOX NO.	NOTES
548	26 555F 00 0F		
549	22 452F 00 0F		
550	00 5F 00 6F		Constants
551	00 6F 00 5F		
552	00 0F 00 0F		
553	5K 2087F 99 2557F		
554	00 0F 00 0F		
555	40 1F K5 5F		Start of drum subroutine
556	42 564F 10 20F	75	Store i, q+1 and y
557	42 2F 32 658F	76	Read or write ?
558	F5 624F 42 568F	77	Store address of basic read order
559	42 589F 42 596F		
560	42 607F L5 623F	78	$\lambda = \lambda'$
561	42 618F 42 610F		
562	L4 555F 42 621F		
563	F4 555F 42 620F		
564	41 0F L5 0F	81	Form addresses for read (write) loops

LOCATION	ORDER	BOX NO.	NOTES	PAGE 25
565	42 OF F5 564F			
566	42 641F 50 1F			
567	75 334F S5 636F			
568	40 3F L5 OF			
569	L0 334F L4 OF			
570	L4 3F 40 625F			
571	40 643F L4 2F			
572	L4 OF 42 626F			
573	42 642F L5 333F			
574	L4 2F 42 613F			
575	42 614F 42 615F			
576	42 616F L5 2F			
577	L4 1F F0 335F			
578	42 611F 42 612F			
579	42 631F 42 647F			
580	42 637F 42 653F			

LOCATION	ORDER	BOX NO.	NOTES
581	L1 1F L4 334F	83	$i \geq 65 ?$
582	32 605F L5 2F	85	Form addresses for read (write) loops
583	L4 334F 42 611F		$\lambda_1 = \lambda_1$
584	42 612F 42 631F		
585	42 647F L5 637F		
586	42 613F 42 614F		
587	L5 620F 42 650F		
588	L5 567F 42 633F		
589	L5 3F L4 0F		
590	L4 1F L0 619F		
591	40 636F 40 654F		
592	L5 0F L0 334F	86	$k > 65$
593	36 603F L5 618F	88	$\lambda_j = \lambda_j$ $r = 65$
594	42 602F L5 617F		
595	46 597F 46 601F		
596	50 334F L5 0F	89	Construct addresses
597	74 0F 85 641F		

LOCATION	ORDER	BOX NO.	NOTES
598	L4 338F L4 1F		
599	L0 667F 40 630F		
600	40 648F L5 2F		
601	L4 OF 42 631F		
602	42 647F 26 OF		
603	L5 621F 42 602F	87	$\lambda_j = \lambda_4, r = k$
604	47 597F 47 601F		
605	26 596F L5 597F	84	$\lambda_1 = \lambda_2$
606	42 633F 42 650F		
607	L5 3F L4 OF		
608	L4 1F F0 334F		
609	40 630F 40 648F		
610	22 610F 26 OF		
611	K2 626F 40 OF		Test constants used in boxes 93, 97, 101, 93', 97', and 101'.
612	K2 642F L5 OF		
613	22 631F 40 OF		
614	22 647F L5 OF		

LOCATION	ORDER	BOX NO.	NOTES	PAGE 28	M 20
615	K2 637F 40 OF				
616	K2 653F L5 OF				
617	00 334F 00 OF		Constants		
618	00 OF 00 OF				
619	00 OF 01 258F				
620	00 OF 00 OF				
621	00 OF 00 OF				
622	00 OF 00 625F				
623	00 OF 00 642F				
624	00 OF 00 336F				
625	00 OF 00 OF	91'			
626	22 626F 40 OF	90'			
627	F5 625F 40 625F	92'			
628	F5 626F 42 626F		Read loops		
629	L0 611F 36 625F	93'			
630	00 OF 00 OF	95'			
631	22 631F 40 OF	94'			

LOCATION	ORDER	BOX NO.	NOTES
632	F5 631F 42 631F	96'	
633	L0 613F 36 OF	97'	
634	L5 630F L4 334F		
635	40 630F 26 630F		
636	00 OF 00 OF	99'	
637	22 637F 40 OF	98'	
638	F5 636F 40 636F	100'	
639	F5 637F 42 637F		
640	L0 615F 36 636F	101'	
641	22 641F 26 OF	102	
642	22 642F L5 OF	90	
643	00 OF 00 OF	91	
644	F5 643F 40 643F	92	Write loops
645	F5 642F 42 642F		
646	L0 612F 32 642F	93	
647	22 647F L5 OF	94	
648	00 OF 00 OF	95	

LOCATION	ORDER	BOX NO.	NOTES	PAGE 30 M 20
649	F5 647F 42 647F	96		
650	L0 614F 32 OF	97		
651	L5 648F L4 334F			
652	40 648F 22 647F			
653	22 653F L5 OF	98		
654	00 OF 00 OF	99		
655	F5 654F 40 654F	100		
656	F5 653F 42 653F			
657	L0 616F 32 653F	101		
658	22 641F L5 624F	79	Store address of basic write order	
659	42 568F 42 589F			
660	42 596F 42 607F			
661	L5 622F 26 561F	80	$\underline{\lambda} = \lambda$	
662	K2 78F 40 936F		Test constant used in read back of Jacobi form	
663	F5 336F F4 335F		Reset same after 130 elements have been read, and adjust test const so this block is disobeyed in future.	
664	40 71F F5 71F			
665	40 77F L5 668F			

LOCATION	ORDER	BOX NO.	NOTES	PAGE 31
666	40 662F			
	26 71F			
667	00 F			
	01 193F			
668	00 F			
	00 F			
	00 800K		Sum check on Routine. If it is to be modified, delete from here on.	
800	81 40F			
	L2 F			
801	40 F			
	L3 F			
802	34 17F			
	FF 54F			
	26 800N			
	58 7543NL4N		Sum check const.	